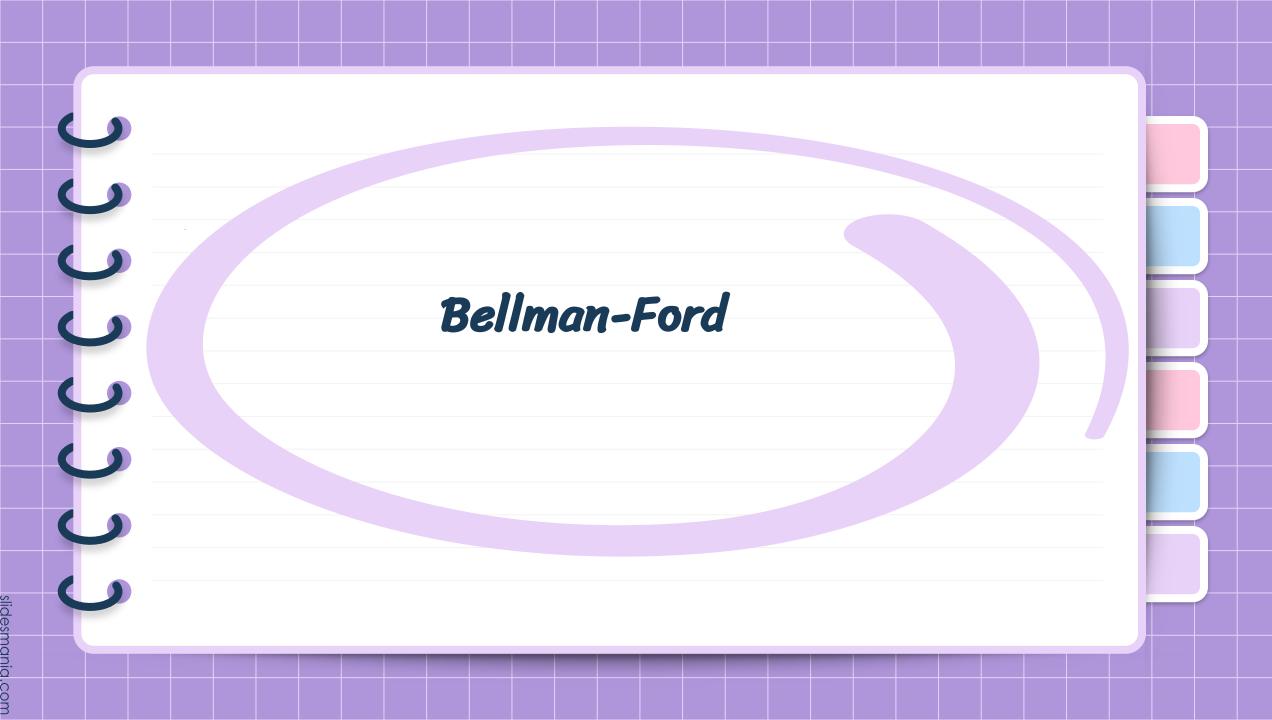
DSC 40B Lecture 25 : Shortest Paths in Weighted Graphs. Bellman-Ford





Bellman-Ford

- Dijkstra's algorithm works only with **positive** weights.
- Bellman-Ford allows us to find shortest distances in the graphs with negative weights as well.

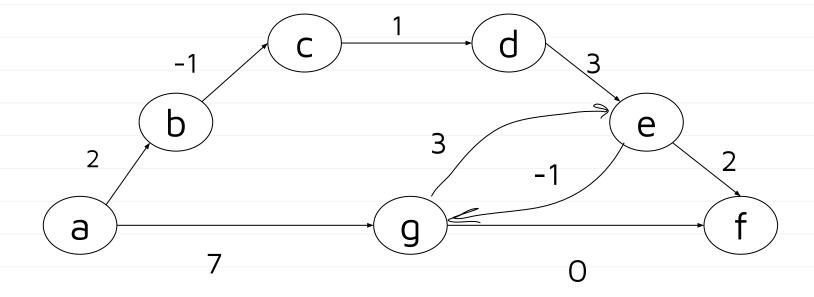


Intuition

- Shortest paths that have many edges are "harder" to discover.
 - May require many updates.
- Shortest paths that have few edges are "easier" to discover.
- Once we've discovered all of the shortest paths with few edges, it makes it easier to discover the shortest paths with more edges.

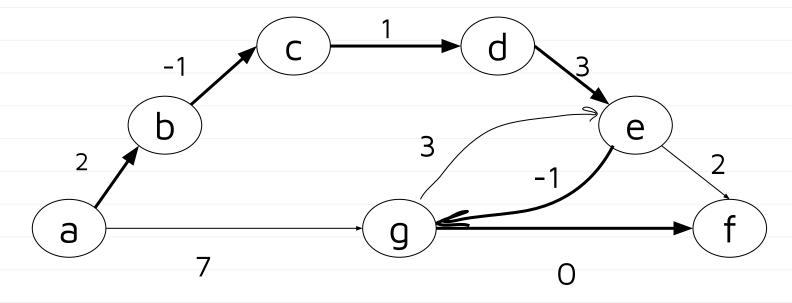
Example

What is the shortest path from **a** to **f**?



Example

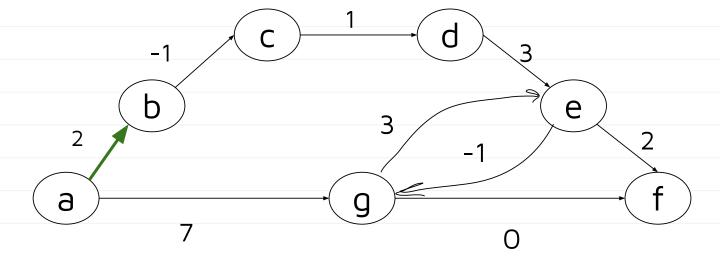
What is the shortest path from **a** to **f**?



Path: a, b, c, d, e, g, f

Updating All Edges

- Suppose we update all of the edges, one by one.
- Then all nodes whose shortest path from s has only one edge are guaranteed to be estimated correctly.





Loop Invariant

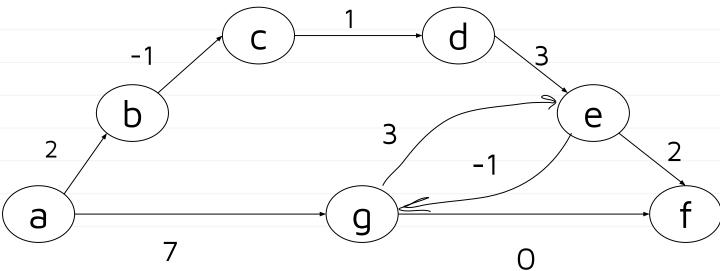
- **One** iteration: update *all* edges in arbitrary order.
- **Loop invariant**: After α iterations, all nodes whose shortest path from s has $\leq \alpha$ edges are **guaranteed** to be estimated correctly.

The Bellman-Ford Algorithm

```
def bellman_ford(graph, weights, source):
"""Assume graph is directed."""
  est = {node: float('inf') for node in graph.nodes}
  est[source] = 0
  predecessor = {node: None for node in graph.nodes}

for i in range(?):
  for (u, v) in graph.edges:
    update(u, v, weights, est, predecessor)
  return est, predecessor
```

Suppose graph.edges returns edges in following order: (c, d), (a, b), (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g)



Suppose graph.edges returns edges in following order: (c, d), (a, b), (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g) ∞ ∞ ∞ ∞

Suppose graph.edges returns edges in following order: (c, d), (a, b), (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g) Do we update distance for d? ∞ ∞ ∞ ∞

Suppose graph.edges returns edges in following order: **(c, d)**, (a, b), (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g) ∞ ∞ ∞

Do we update distance for d? **No**

 ∞

Suppose graph.edges returns edges in following order: (c, d), (a, b), (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g) Do we update distance for b? ∞ ∞ ∞ ∞

Suppose graph.edges returns edges in following order: **(c, d)**, **(a, b)**, (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g) ∞ ∞ ∞

Do we update distance for b? **Yes**

 ∞

Suppose graph.edges returns edges in following order: (c, d), (a, b), (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g) Do we update distance for c? ∞ ∞ ∞ ∞

Suppose graph.edges returns edges in following order: **(c, d)**, **(a, b)**, **(b, c)**, (g, f), (e, g), (g, e), (d, e), (e, f), (a, g) Do we update distance for c? Yes ∞ ∞ ∞

Suppose graph.edges returns edges in following order: (c, d), (a, b), (b, c), **(g, f)**, (e, g), (g, e), (d, e), (e, f), (a, g) Do we update distance for f? No ∞ ∞ ∞

Suppose graph.edges returns edges in following order: (c, d), (a, b), (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g) Do we update distance for g? No ∞ ∞ ∞

Suppose graph.edges returns edges in following order: (c, d), (a, b), (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g) Do we update distance for e? No ∞ ∞ ∞

Suppose graph.edges returns edges in following order: (c, d), (a, b), (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g) Do we update distance for e? No ∞ ∞ ∞

Suppose graph.edges returns edges in following order: (c, d), (a, b), (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g) Do we update distance for f? No ∞ ∞ ∞

Suppose graph.edges returns edges in following order: (c, d), (a, b), (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g) Do we update distance for g? ∞ ∞ ∞

Suppose graph.edges returns edges in following order: (c, d), (a, b), (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g) Do we update distance for g? **Yes** ∞ ∞ ∞

Suppose graph.edges returns edges in following order: (c, d), (a, b), (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g) It was just a **single** iteration of BF ∞ ∞

Suppose graph.edges returns edges in following order: (c, d), (a, b), (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g) We only can guarantee one correct edge after this iteration: (a, b) ∞ ∞

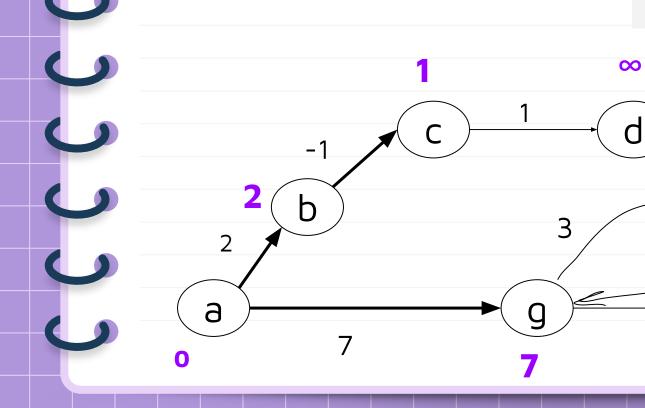
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Suppose graph.edges returns edges in following order:

(c, d), (a, b), (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g)

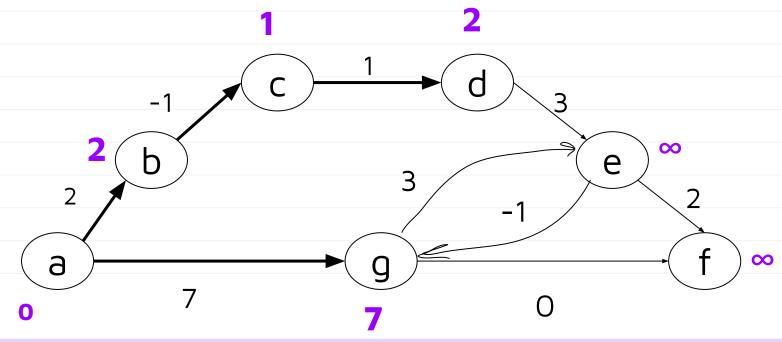
Do we update distance for d? **Yes**

 ∞



(c, d), **(a, b)**, (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g)

Do we update distance for b? No



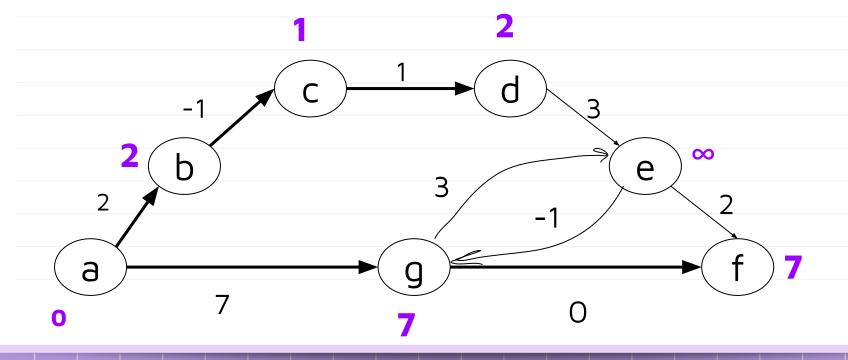
Suppose graph.edges returns edges in following order: **(c, d)**, **(a, b)**, **(b, c)**, (g, f), (e, g), Do we update distance for c? No (g, e), (d, e), (e, f), (a, g) ∞

Suppose graph.edges returns edges in following order: (c, d), (a, b), (b, c), **(g, f)**, (e, g), Do we update distance for f? **Yes** (g, e), (d, e), (e, f), (a, g) ∞

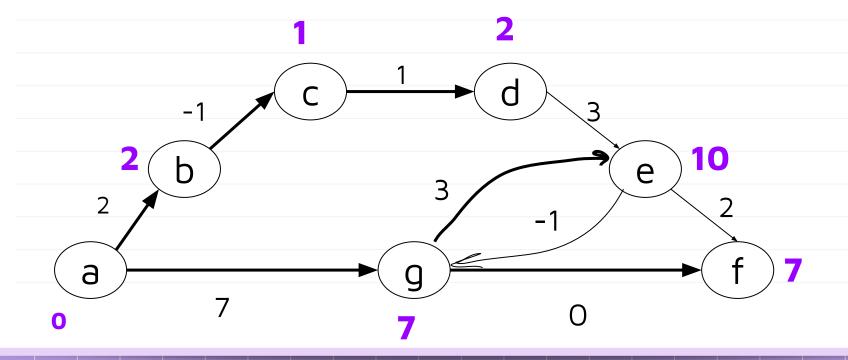
Suppose graph.edges returns edges in following order: (c, d), (a, b), (b, c), **(g, f)**, (e, g), Do we update distance for f? **Yes** (g, e), (d, e), (e, f), (a, g)

Suppose graph.edges returns edges in following order: (c, d), (a, b), (b, c), (g, f), (e, g), Do we update distance for g? No (g, e), (d, e), (e, f), (a, g)

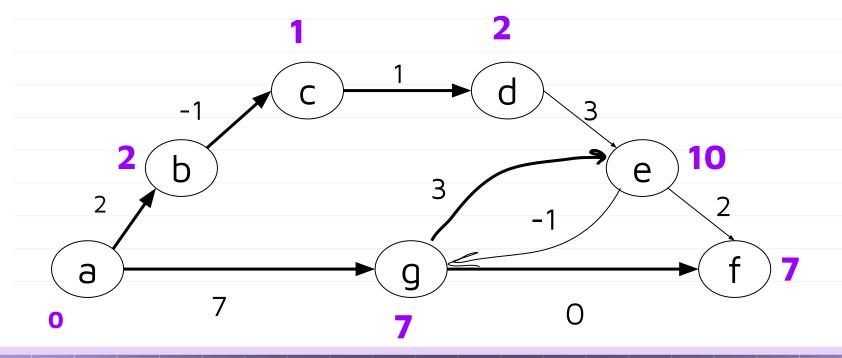
Do we update distance for e? **Yes**



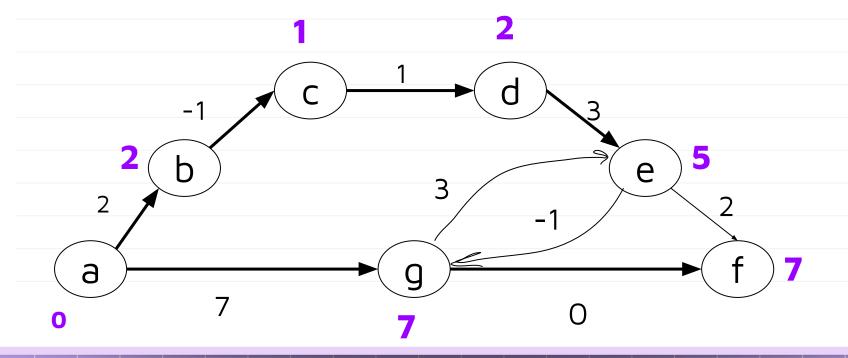
Do we update distance for e? **Yes**



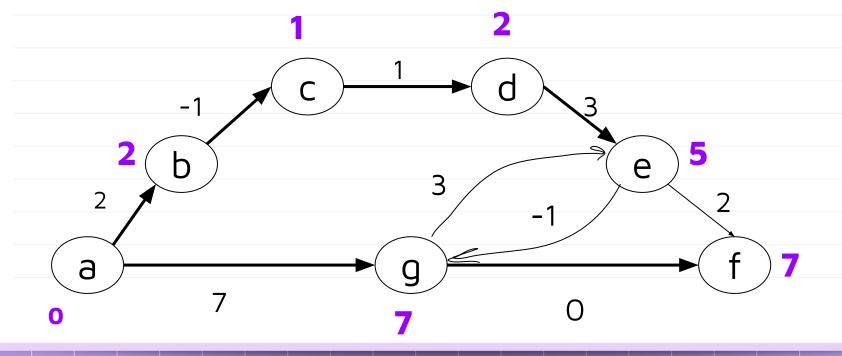
Do we update distance for e? **Yes**



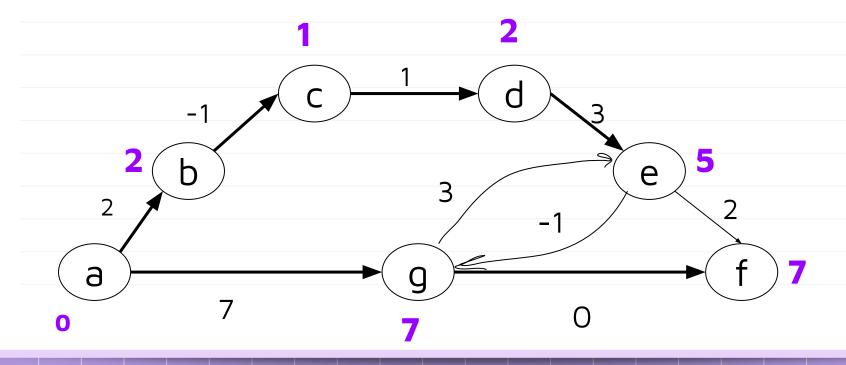
Do we update distance for f? No



Do we update distance for g? No



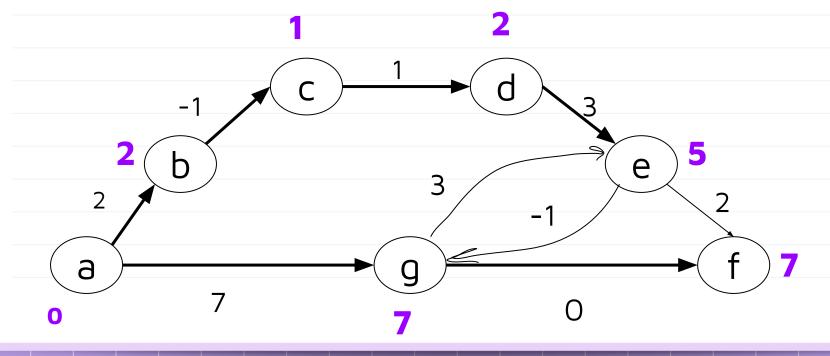
It was the end of the **second** iteration of BF



(c, d), (a, b), (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g)

One more time



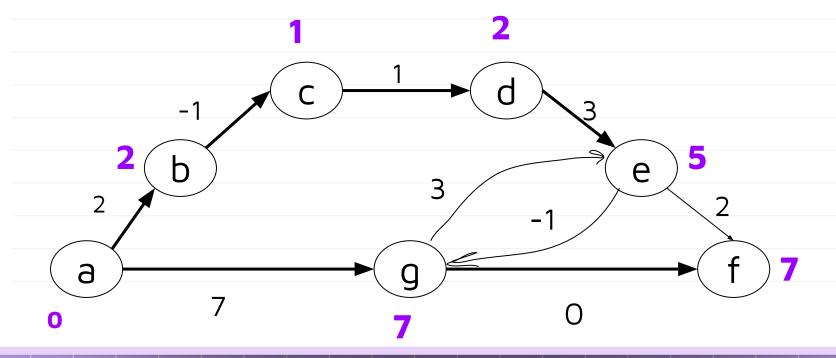


Suppose graph.edges returns edges in following order:

(c, d), (a, b), (b, c), (a, f), (e, a),

(c, d), (a, b), (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g)

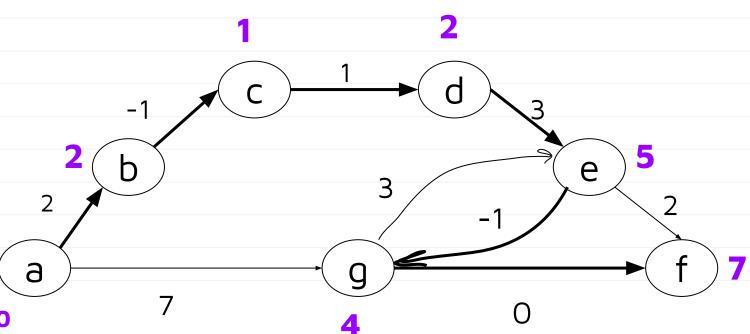
No change



Suppose graph.edges returns edges in following order:

(c, d), (a, b), (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g)

It was the end of the third iteration of BF





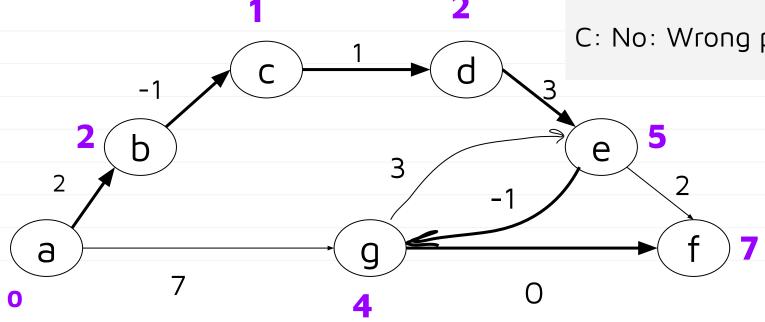
(c, d), (a, b), (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g)

Is everything look correct?

A: Yes, we are done!

B: No: Wrong number

C: No: Wrong predecessor





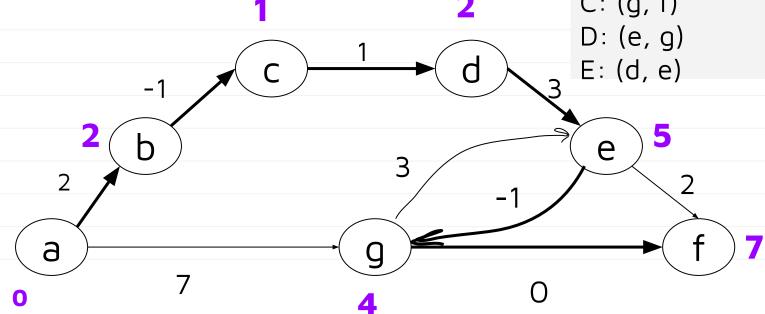
(c, d), (a, b), (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g)

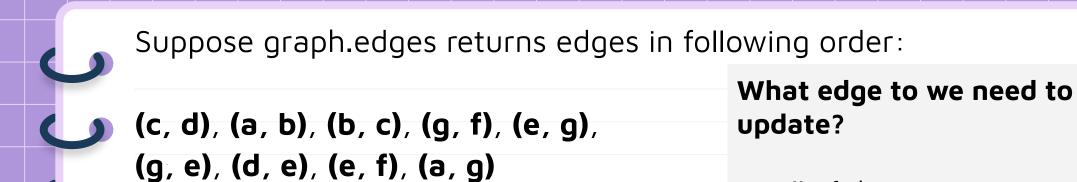
What edge to we need to update?

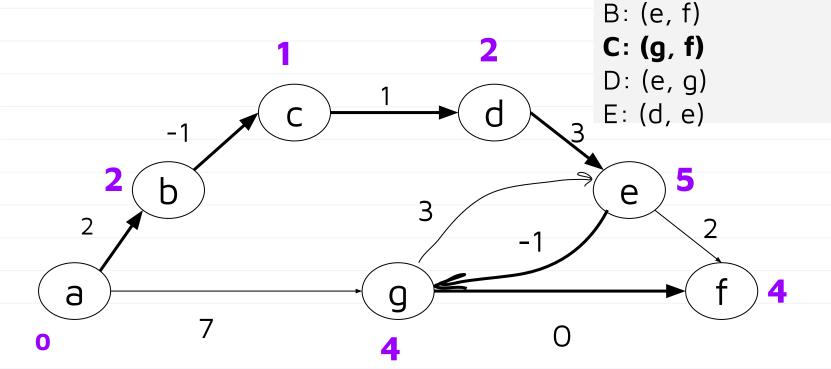
A: All of them

B: (e, f)

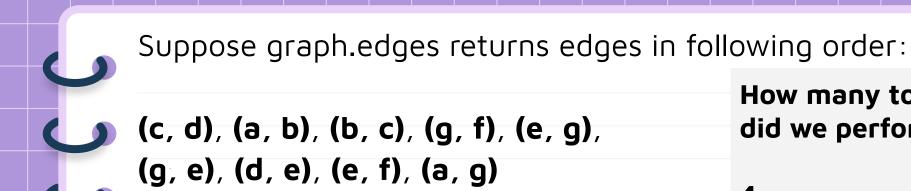
C: (g, f)





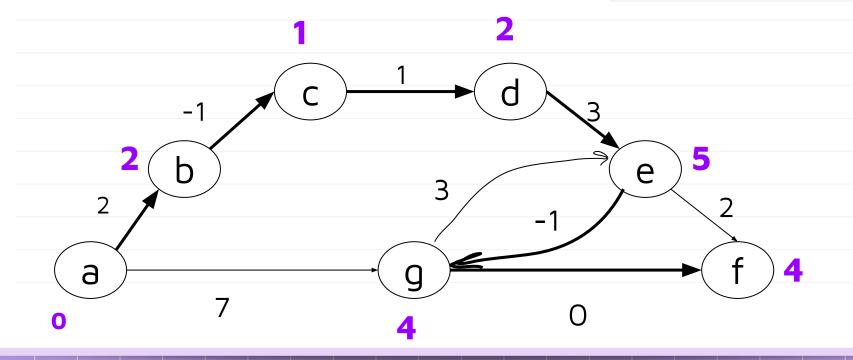


A: All of them



How many total iterations did we perform?

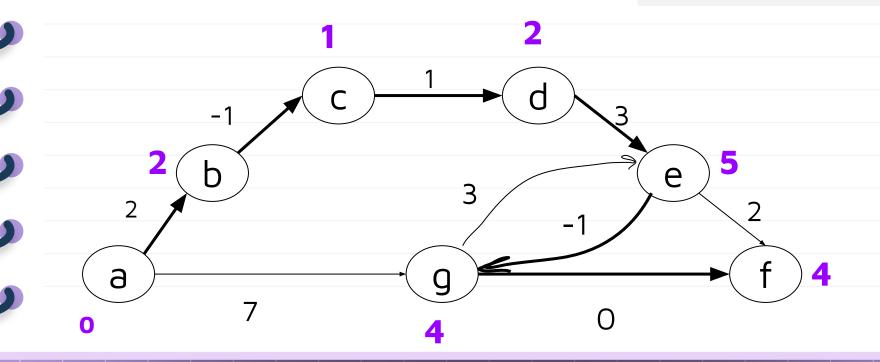
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(c, d), (a, b), (b, c), (g, f), (e, g), (g, e), (d, e), (e, f), (a, g) Do we have to run one more iteration?

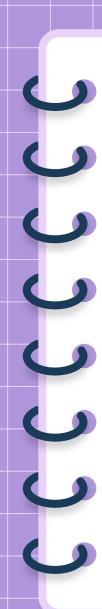
No





Loop Invariant

- **One** iteration: update *all* edges in arbitrary order.
- **Loop invariant**: After α iterations, all nodes whose **actual** shortest path from s has $\leq \alpha$ edges are guaranteed to be estimated **correctly**.



Bellman-Ford

- Claim: each node must have a shortest path which is simple (Edge case: cycles of weight zero)
- The most edges a simple path can have is |V| 1
- **Idea of Bellman-Ford**: iteratively update all edges, repeat |V| 1 times.

The Bellman-Ford Algorithm

```
def bellman_ford(graph, weights, source):
"""Assume graph is directed."""
  est = {node: float('inf') for node in graph.nodes}
  est[source] = 0
  predecessor = {node: None for node in graph.nodes}

for i in range(len(graph.nodes)-1):
  for (u, v) in graph.edges:
     update(u, v, weights, est, predecessor)
  return est, predecessor
```

```
The Bellman-Ford Algorithm: Complexity
 def bellman_ford(graph, weights, source):
 """Assume graph is directed."""
    est = {node: float('inf') for node in graph.nodes}
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    predecessor = {node: None for node in graph.nodes}
    for i in range(len(graph.nodes)-1):
        for (u, v) in graph.edges:
           update(u, v, weights, est, predecessor)
    return est, predecessor
 Setup: _____
                               time:
 Each update takes _____
                               _time
 There are exactly ____updates
 Total time complexity:
```

```
The Bellman-Ford Algorithm: Complexity
 def bellman_ford(graph, weights, source):
 """Assume graph is directed."""
    est = {node: float('inf') for node in graph.nodes}
    est[source] = 0
    predecessor = {node: None for node in graph.nodes}
    for i in range(len(graph.nodes)-1):
        for (u, v) in graph.edges:
           update(u, v, weights, est, predecessor)
    return est, predecessor
 Setup: O (V)
                                time:
 Each update takes _____
                               time
 There are exactly ____updates
 Total time complexity:_____
```

```
The Bellman-Ford Algorithm: Complexity
 def bellman_ford(graph, weights, source):
 """Assume graph is directed."""
    est = {node: float('inf') for node in graph.nodes}
    est[source] = 0
    predecessor = {node: None for node in graph.nodes}
     for i in range(len(graph.nodes)-1):
        for (u, v) in graph.edges:
            update(u, v, weights, est, predecessor)
     return est, predecessor
 Setup: Θ (V)
                                 time
 Each update takes \underline{\Theta} (1)
                                 _time
 There are exactly ____updates
 Total time complexity:_____
```

```
The Bellman-Ford Algorithm: Complexity
 def bellman_ford(graph, weights, source):
 """Assume graph is directed."""
     est = {node: float('inf') for node in graph.nodes}
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     predecessor = {node: None for node in graph.nodes}
     for i in range(len(graph.nodes)-1):
        for (u, v) in graph.edges:
            update(u, v, weights, est, predecessor)
     return est, predecessor
 Setup: O (V)
                                   time
 Each update takes \underline{\Theta} (1)
                                  time
 There are exactly (|V|-1) \times |E|
                                  _updates
 Total time complexity:
```

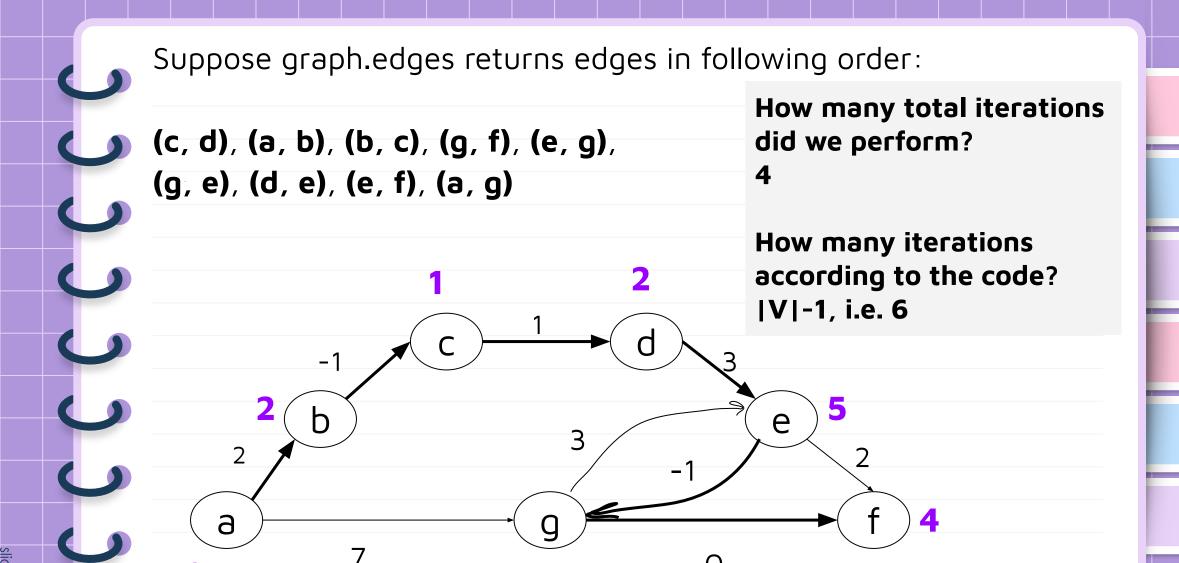
```
The Bellman-Ford Algorithm: Complexity
 def bellman_ford(graph, weights, source):
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     est[source] = 0
     predecessor = {node: None for node in graph.nodes}
     for i in range(len(graph.nodes)-1):
         for (u, v) in graph.edges:
            update(u, v, weights, est, predecessor)
     return est, predecessor
 Setup: Θ (V)
                                   time
 Each update takes \underline{\Theta} (1)
                                   time
 There are exactly (|V|-1) \times |E|
                                   _updates
 Total time complexity: \Theta (V + VE)
```

Early Stopping and Negative Cycles



Early Stopping

- \bullet B-F may not need to run for |V| 1 iterations.
- If no predecessors change, we can break:



Early Stopping

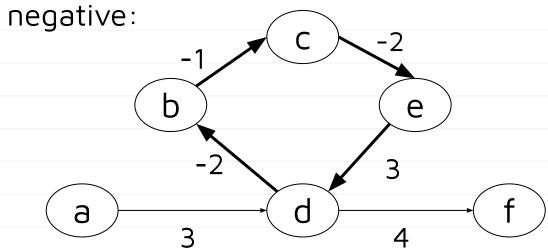
- \bullet B-F may not need to run for |V| 1 iterations.
- If no predecessors change, we can break:

```
def bellman_ford(graph, weights, source):
    """Early stopping version."""
    est = {node: float('inf') for node in graph.nodes}
    est[source] = 0
    predecessor = {node: None for node in graph.nodes}

for i in range(len(graph.nodes) - 1):
    any_changes = False
    for (u, v) in graph.edges:
        changed = update(u, v, weights, est, predecessor)
        any_changes = changed or any_changes
    if not any_changes:
        break
    return est, predecessor
```

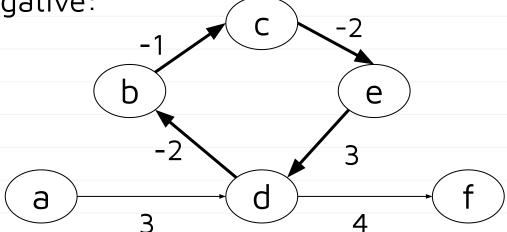
Negative Cycles

• A **negative cycle** is a cycle whose total edge weight is



Negative Cycles

• A **negative cycle** is a cycle whose total edge weight is negative:



 If a graph has a negative cycle, (some) shortest paths are not well defined.



Detecting Negative Cycles

- If graph **does not have** negative cycles, estimated distances eventually stop changing (after at most |V| 1 iterations).
- If graph has negative cycles, estimated distances always decrease.
- To detect them: run a |V| [th iteration; if distances change, a negative cycle exists.

Detecting Negative Cycles

```
def bellman ford(graph, weights, source):
    """Early stopping version, detects negative cycles."""
    est = {node: float('inf') for node in graph.nodes}
    est[source] = 0
    predecessor = {node: None for node in graph.nodes}
    for i in range(len(graph.nodes)):
        any_changes = False
        for (u, v) in graph.edges:
            changed = update(u, v, weights, est, predecessor)
            any_changes = changed or any_changes
        if not any_changes:
            break
    # this will be True if negative cycles exist
    contains_negative_cycles = any_changes
    return est, predecessor, contains_negative_cycles
```

Thank you!

Do you have any questions?

CampusWire!