

DSC 40B

***Lecture 18 : Graph
representations***

Adjacency Matrices

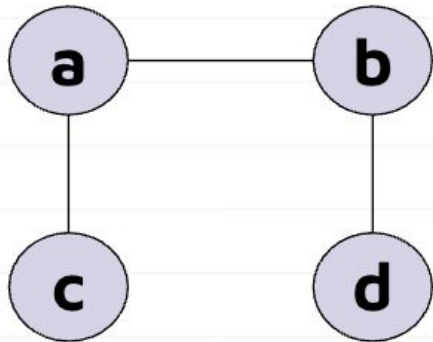
Representations

- How do we store a graph in a computer's memory?
- Three approaches:
 - Adjacency matrices.
 - Adjacency lists.
 - "Dictionary of sets"

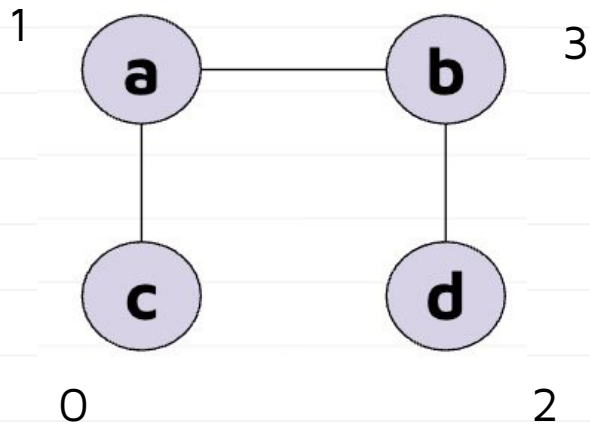
Adjacency Matrices

- Assume nodes are numbered $0, 1, \dots, |V| - 1$
- Allocate a $|V| \times |V|$ (Numpy) array
- Fill array as follows:
 - $\text{arr}[i,j] = 1$ if $(i, j) \in E$
 - $\text{arr}[i,j] = 0$ if $(i, j) \notin E$

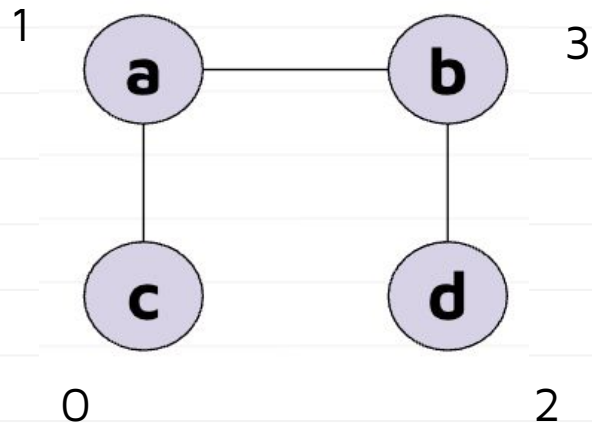
Example



Example

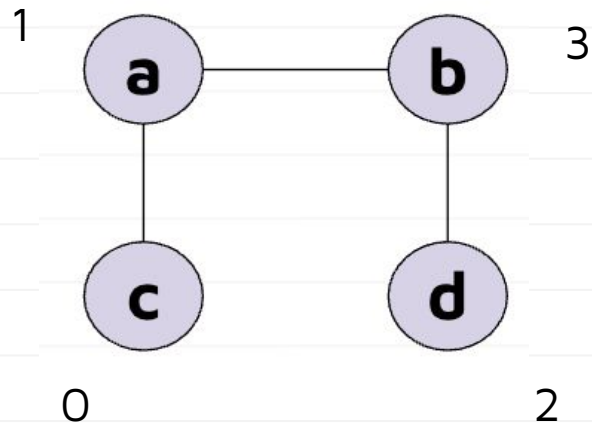


Example



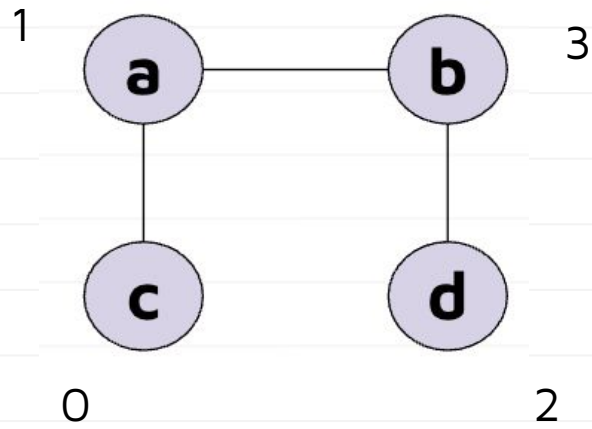
	0	1	2	3
0				
1				
2				
3				

Example



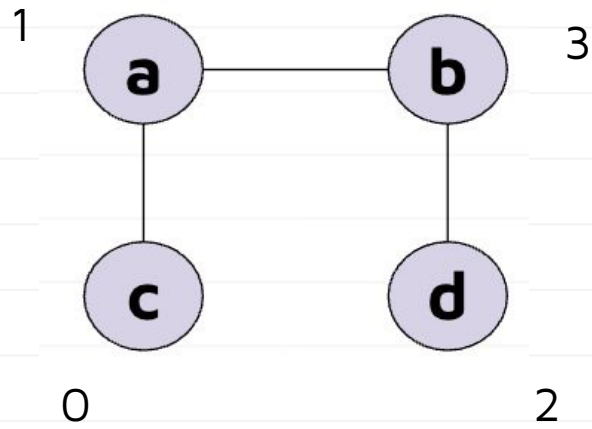
	0	1	2	3
0		1		
1				
2				
3				

Example



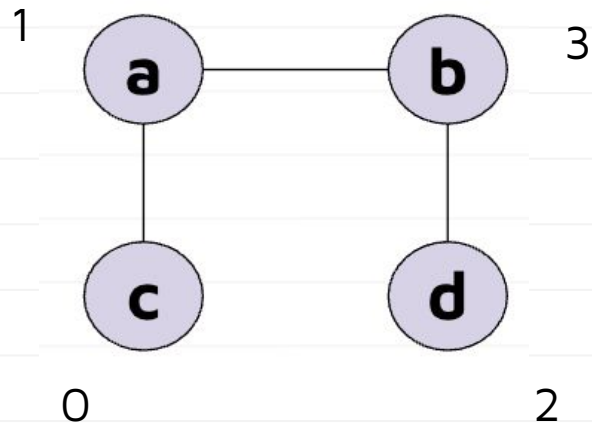
	0	1	2	3
0		1		
1	1			
2				
3				

Example



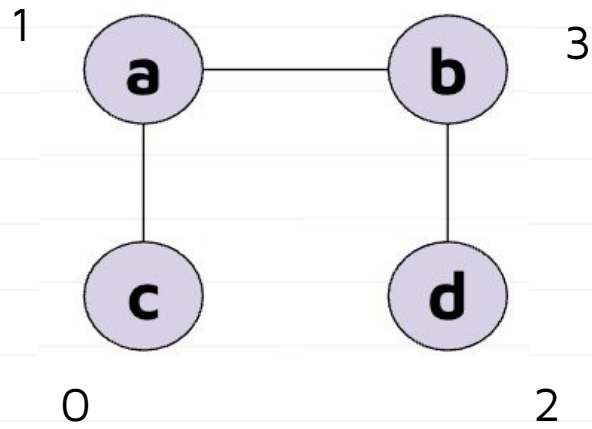
	0	1	2	3
0		1		
1	1			1
2				
3				

Example



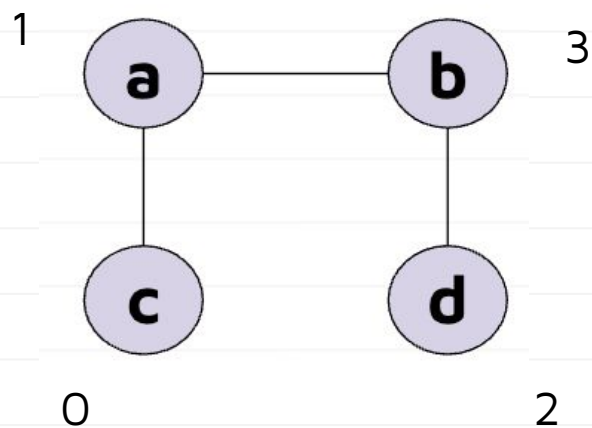
	0	1	2	3
0		1		
1	1			1
2				
3		1		

Example



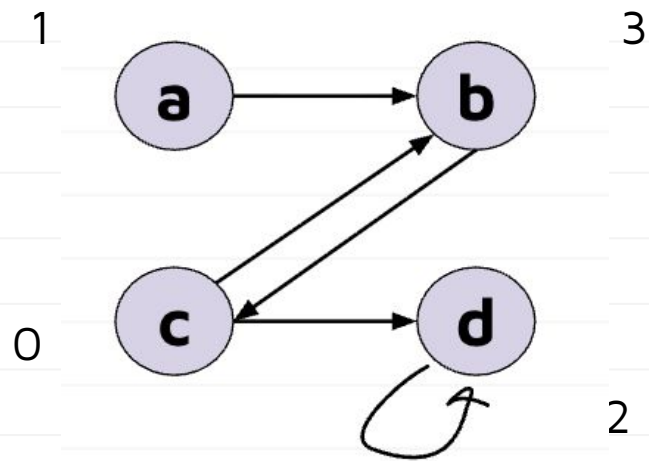
	0	1	2	3
0		1		
1	1			1
2				1
3		1	1	

Example



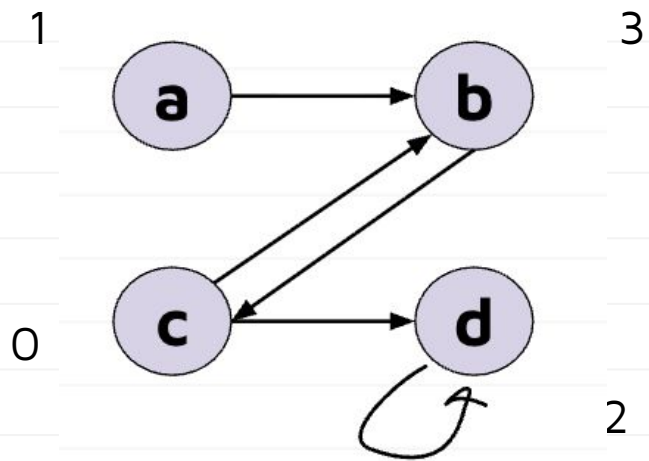
	0	1	2	3
0	0	1	0	0
1	1	0	0	1
2	0	0	0	1
3	0	1	1	0

Example



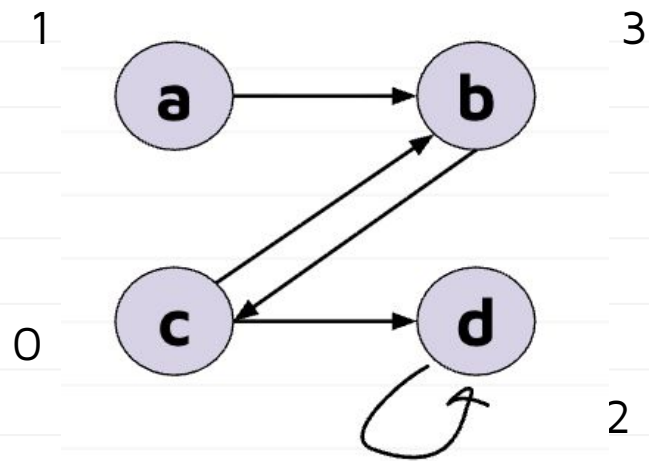
	0	1	2	3
0				
1				1
2				
3				

Example



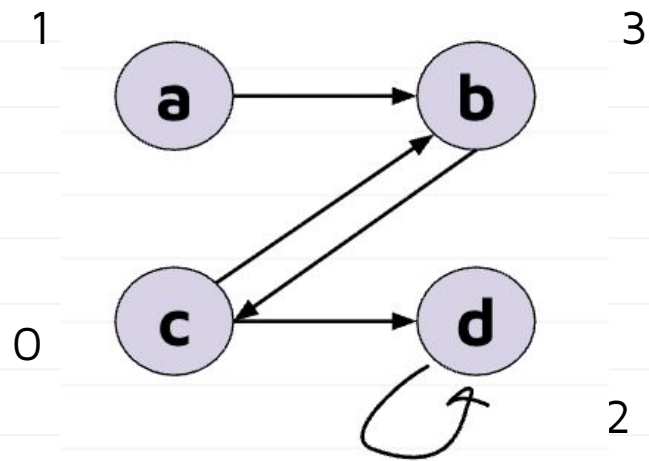
	0	1	2	3
0				1
1				1
2				
3				

Example



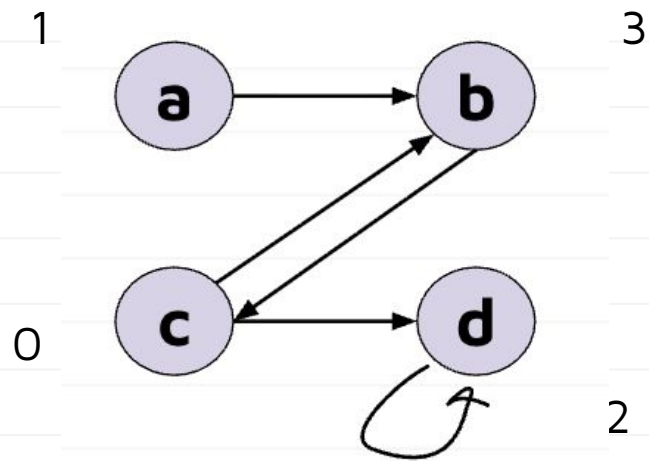
	0	1	2	3
0			1	1
1				1
2				
3				

Example



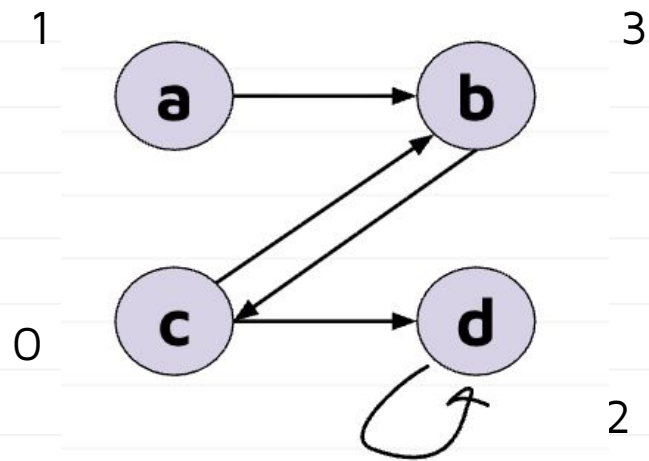
	0	1	2	3
0			1	1
1				1
2				
3	1			

Example



	0	1	2	3
0			1	1
1				1
2			1	
3	1			

Example



	0	1	2	3
0	0	0	1	1
1	0	0	0	1
2	0	0	1	0
3	1	0	0	0

Observations

- If G is **undirected**, matrix is *symmetric*.
- If G is **directed**, matrix *may not* be symmetric.

Time Complexity

operation	code	time
Edge query	<code>adj[i, j] == 1</code>	$\Theta(1)$
Degree(i)	<code>np.sum(adj[i,:])</code>	$\Theta(V)$

Space Requirements

- Uses $|V|^2$ bits, even if there are very few edges.
- But most real-world graphs are **sparse**.
 - They contain many *fewer* edges than possible.

Example: Facebook

- Facebook has 2 billion users. (V)

$$(2 \times 10^9)^2 = 4 \times 10^{18} \text{ bits}$$

= 500 petabits

≈ 6500 years of video at 1080p

≈ 60 copies of the internet as it was in 2000

Adjacency Matrices and Math

- Adjacency matrices are useful mathematically.
- **Example:** (i, j) entry of A^2 gives number of hops of length 2 between i and j .

Adjacency Lists

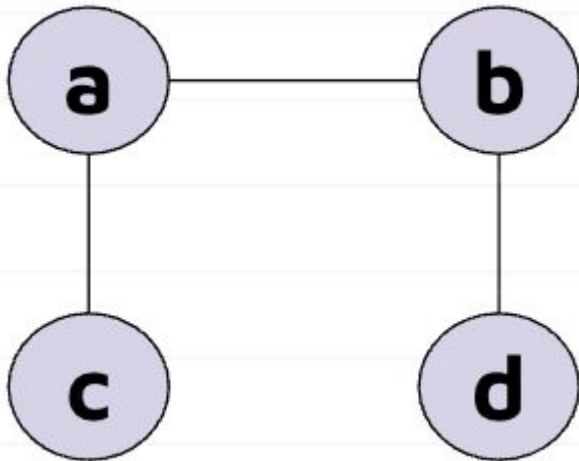
What's Wrong with Adjacency Matrices?

- Requires $\Theta(|V|^2)$ storage.
- Even if the graph has no edges.
- **Idea**: only store the edges that exist.

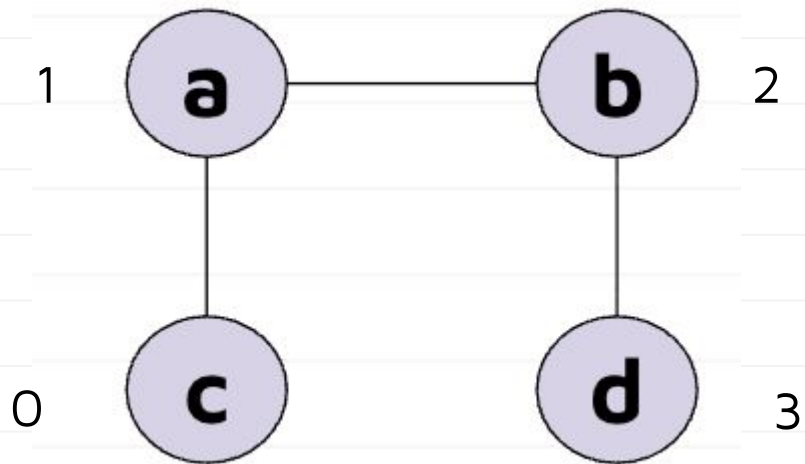
Adjacency Lists

- Create a list `adj` containing $|V|$ lists.
- `adj[i]` is list containing the neighbors of node i .

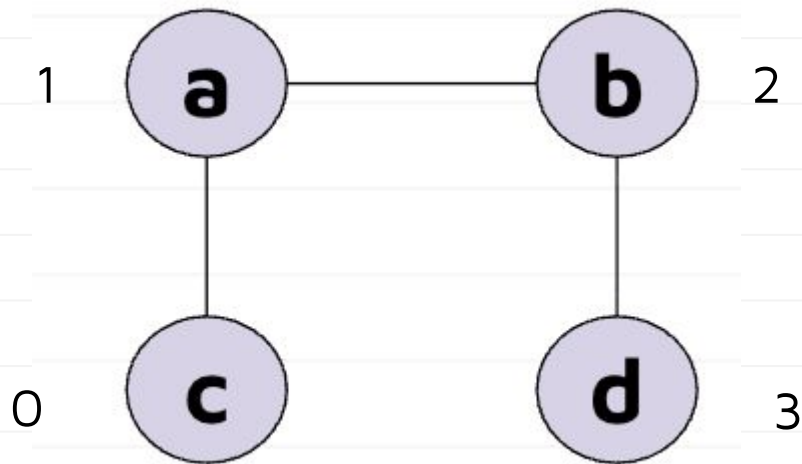
Example



Example



Example



adj = [

[,

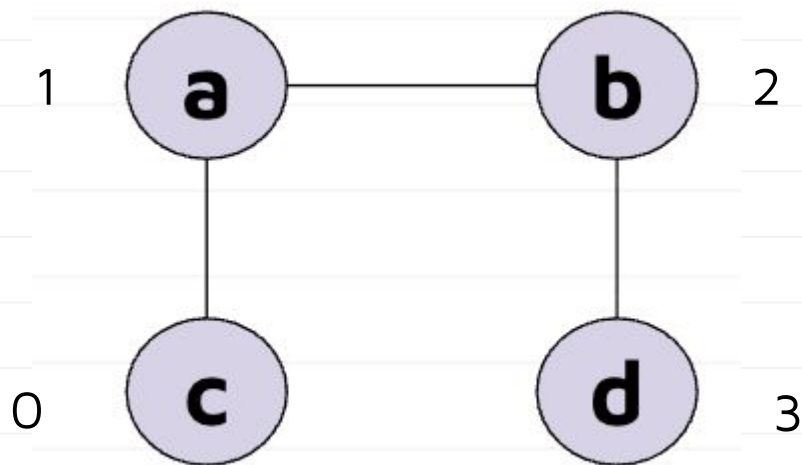
[,

[,

[]

]

Example



adj = [

#0 [,

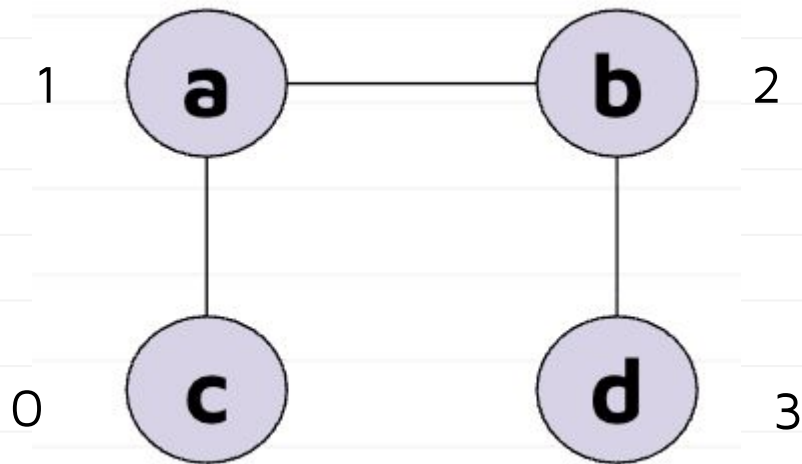
#1 [,

#2 [,

#3 []

]

`adj[i]` is list containing the *neighbors* of node *i*.



```
adj = [
```

```
#0 [    ?    ],
```

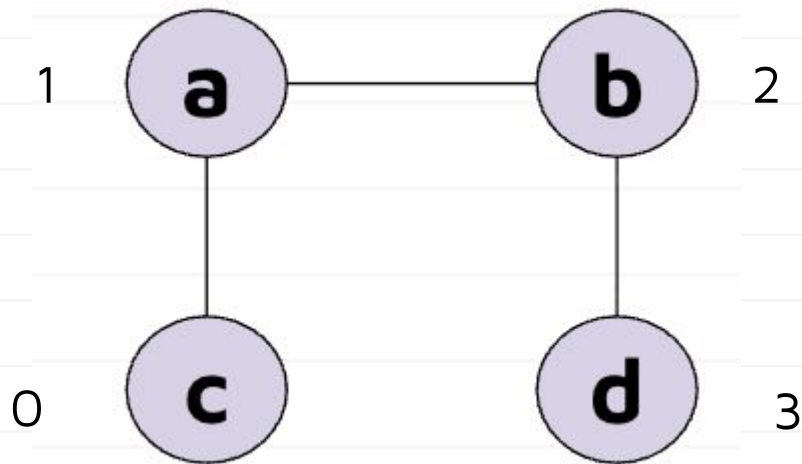
```
#1 [    ?    ],
```

```
#2 [    ?    ],
```

```
#3 [    ?    ]
```

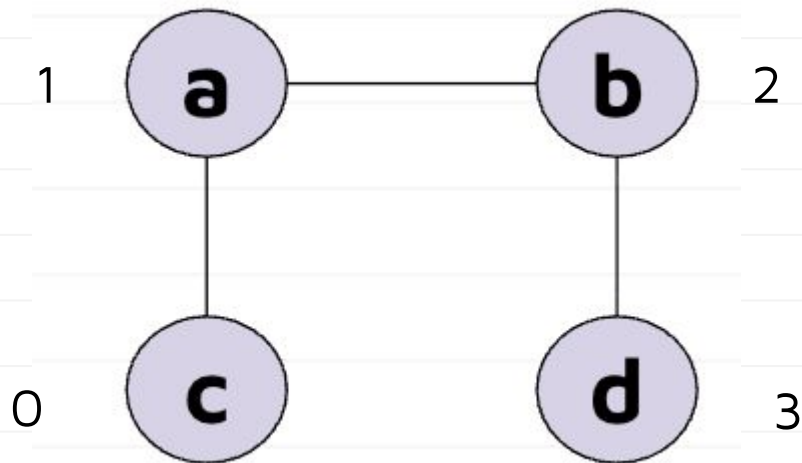
```
]
```

`adj[i]` is list containing the *neighbors* of node *i*.



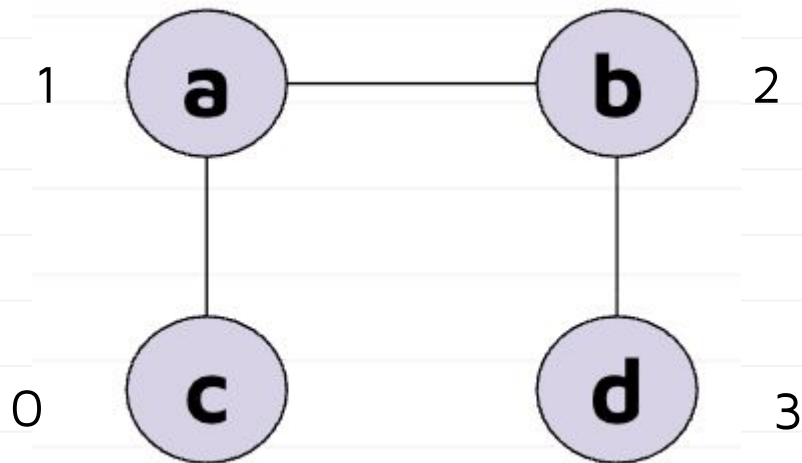
```
adj = [  
    #0 [ 1 ],  
    #1 [    ?    ],  
    #2 [    ?    ],  
    #3 [    ?    ]  
]
```

`adj[i]` is list containing the *neighbors* of node *i*.



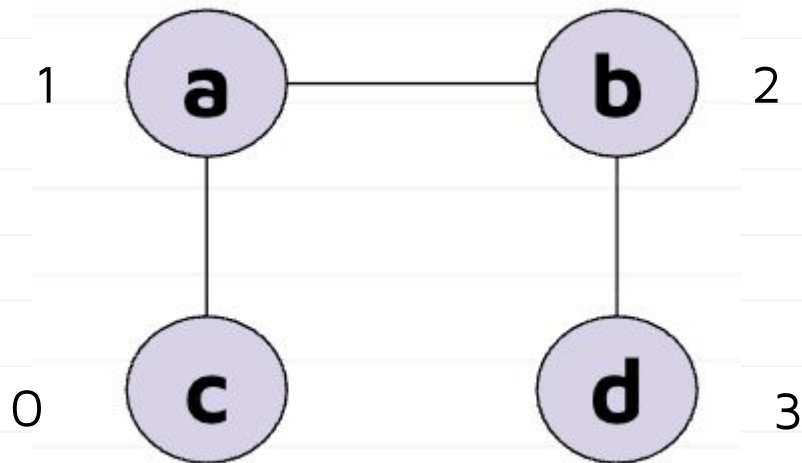
```
adj = [  
    #0 [ 1 ],  
    #1 [ 0, 2 ],  
    #2 [      ?      ],  
    #3 [      ?      ]  
]
```

`adj[i]` is list containing the *neighbors* of node *i*.



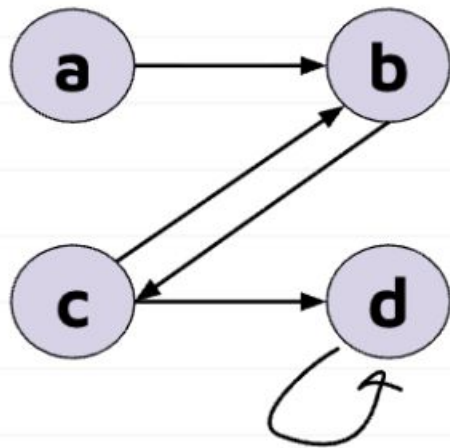
```
adj = [  
    #0 [ 1 ],  
    #1 [ 0, 2 ],  
    #2 [ 1, 3 ],  
    #3 [      ?      ]  
]
```

`adj[i]` is list containing the *neighbors* of node *i*.

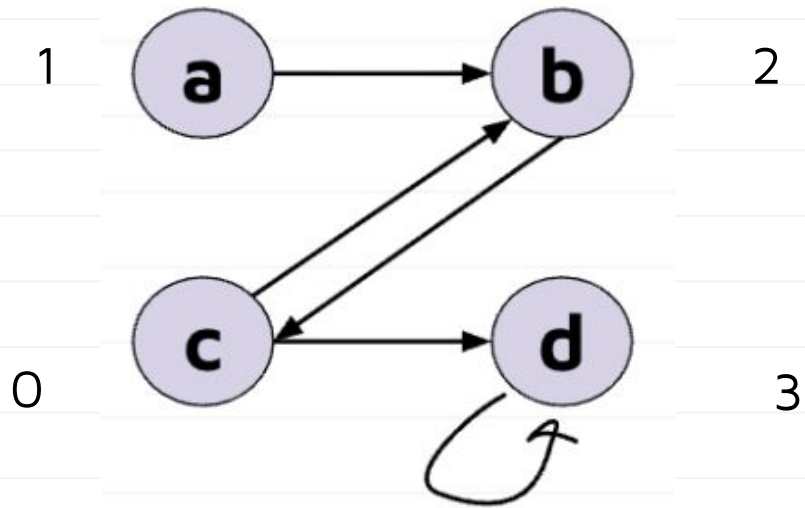


```
adj = [  
    #0 [ 1 ],  
    #1 [ 0, 2 ],  
    #2 [ 1, 3 ],  
    #3 [ 2 ]  
]
```

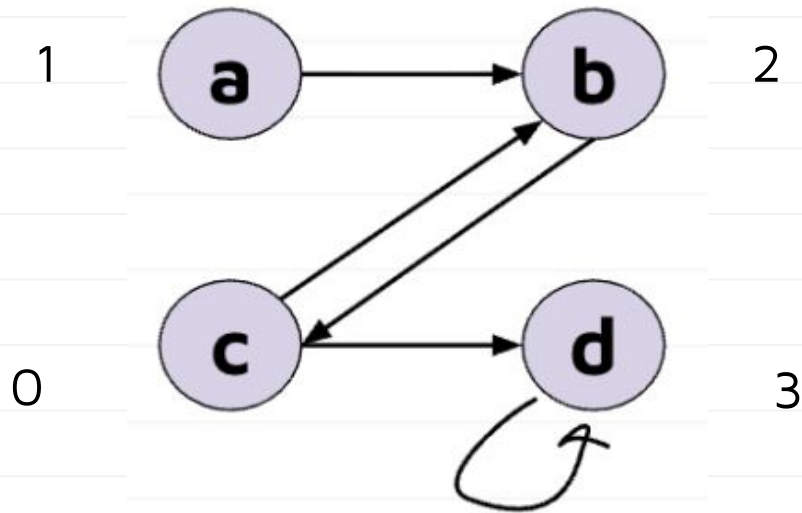
Example



Example



Example



adj = [

[,

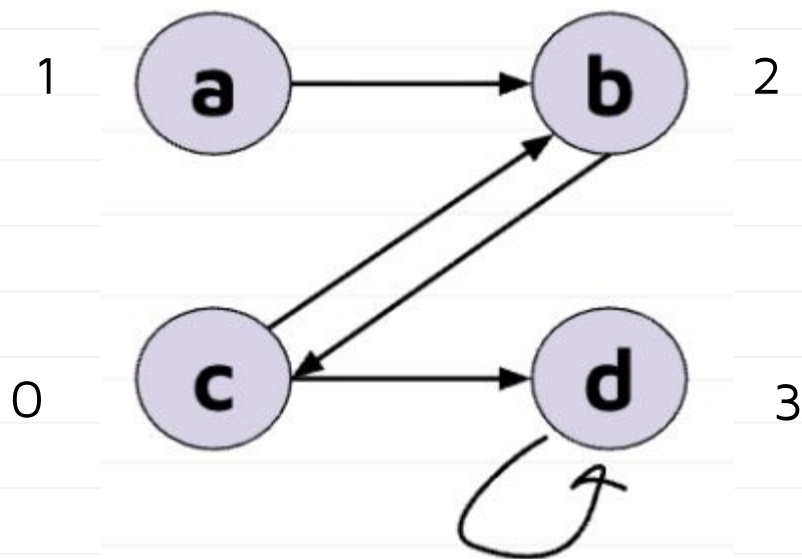
[,

[,

[]

]

Example



```
adj = [
```

```
  #0 [
```

```
    ],
```

```
  #1 [
```

```
    ],
```

```
  #2 [
```

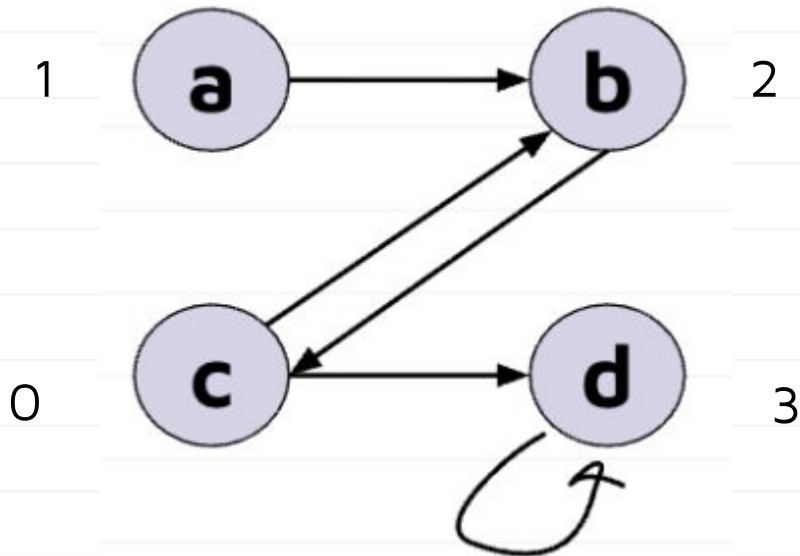
```
    ],
```

```
  #3 [
```

```
    ]
```

```
]
```

`adj[i]` is list containing the *neighbors* of node *i*.



```
adj = [
```

```
#0 [    ?    ],
```

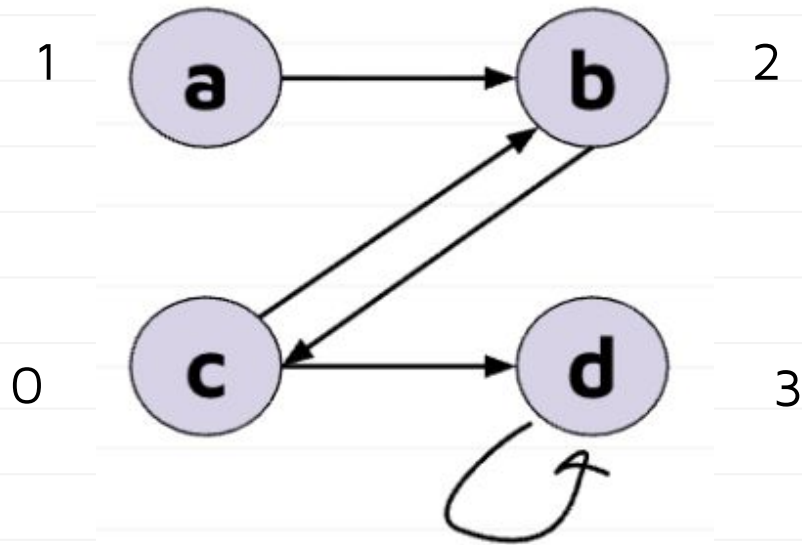
```
#1 [    ?    ],
```

```
#2 [    ?    ],
```

```
#3 [    ?    ]
```

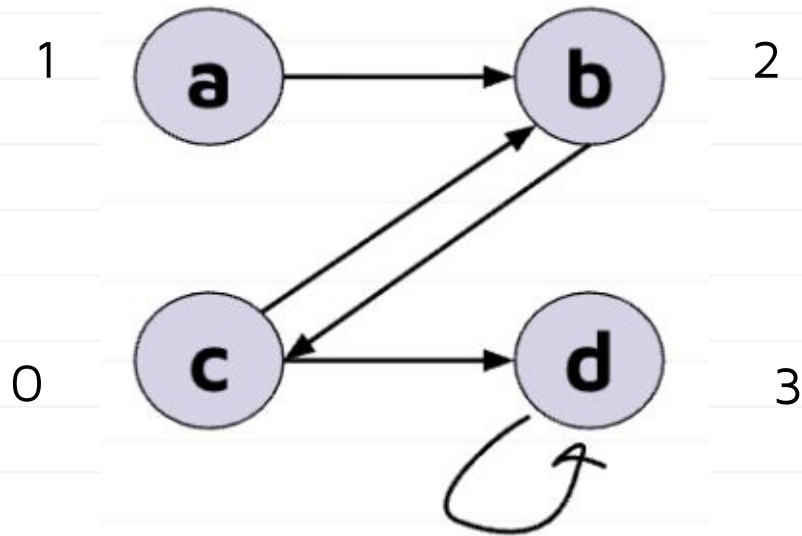
```
]
```

`adj[i]` is list containing the *neighbors* of node *i*.



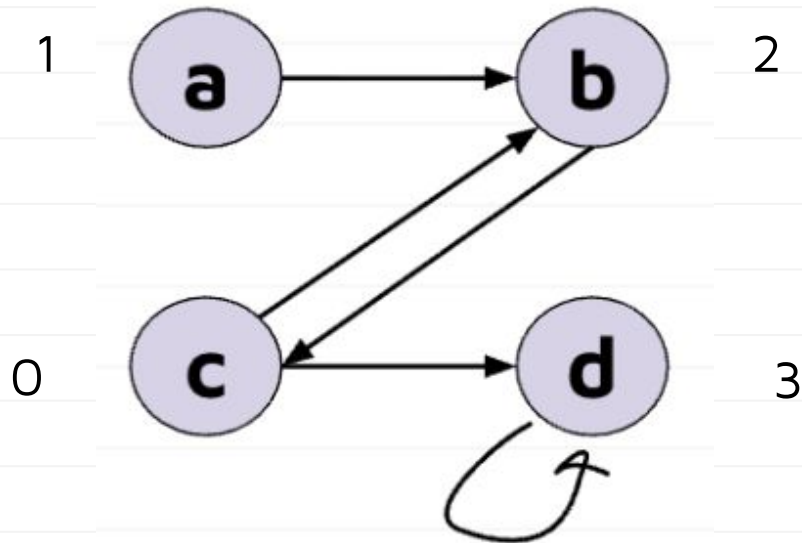
```
adj = [  
    #0 [ 2, 3 ],  
    #1 [      ?      ],  
    #2 [      ?      ],  
    #3 [      ?      ]  
]
```

`adj[i]` is list containing the *neighbors* of node *i*.



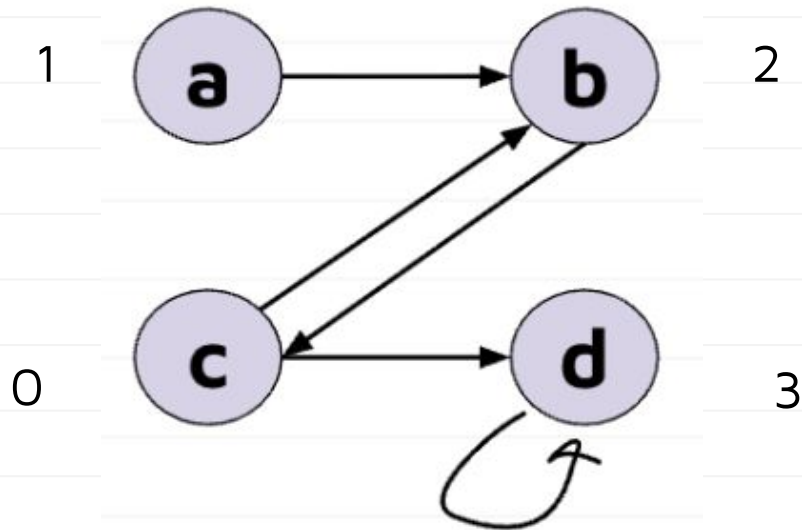
```
adj = [  
    #0 [ 2, 3 ],  
    #1 [ 2 ],  
    #2 [      ?      ],  
    #3 [      ?      ]  
]
```

`adj[i]` is list containing the *neighbors* of node *i*.



```
adj = [  
    #0 [ 2, 3 ],  
    #1 [ 2 ],  
    #2 [ 0 ],  
    #3 [      ?      ]  
]
```

`adj[i]` is list containing the *neighbors* of node *i*.



```
adj = [  
    #0 [ 2 ],  
    #1 [ 2 ],  
    #2 [ 0 ],  
    #3 [ 3 ]  
]
```

Observations

- If G is **undirected**, each edge appears **twice**.
- If G is **directed**, each edge appears **once**.

Time Complexity

operation	code	time
Edge query	<code>j in adj[i, j]</code>	$\Theta(\text{degree}(i))$
Degree(i)	<code>len(adj[i])</code>	$\Theta(1)$

Exercise 1: Build a graph from the list

```
adj = [  
    [1],  
    [0, 2],  
    [1, 3],  
    [2]  
]  
  
for u in range(len(adj)):  
    for v in adj[u]:  
        print(f"({u}, {v})")
```

Exercise 1: Build a graph from the list

0

1

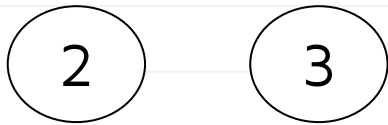
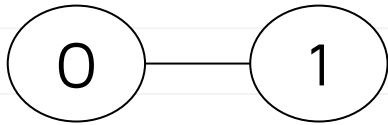
2

3

```
adj = [  
    [1],  
    [0, 2],  
    [1, 3],  
    [2]  
]
```

```
for u in range(len(adj)):  
    for v in adj[u]:  
        print(f"({u}, {v})")
```

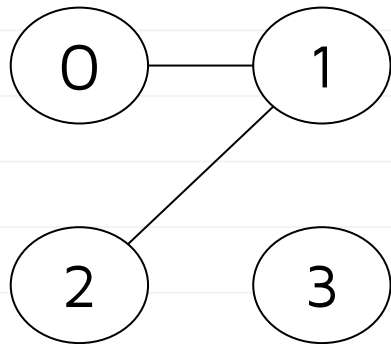
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  [1, 3],  
  [2]  
]
```

```
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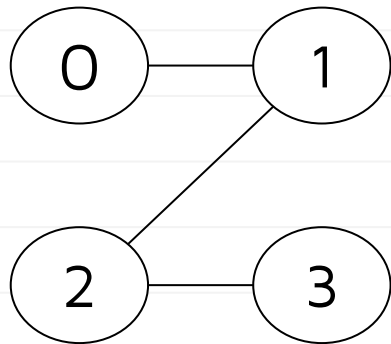
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```
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    [1, 3],  
    [2]  
]
```

```
for u in range(len(adj)):  
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        print(f"({u}, {v})")
```

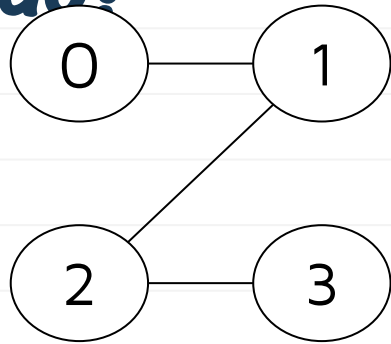
Exercise 1: Build a graph from the list



```
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    [1],  
    [0, 2],  
    [1, 3],  
    [2]  
]
```

```
for u in range(len(adj)):  
    for v in adj[u]:  
        print(f"({u}, {v})")
```

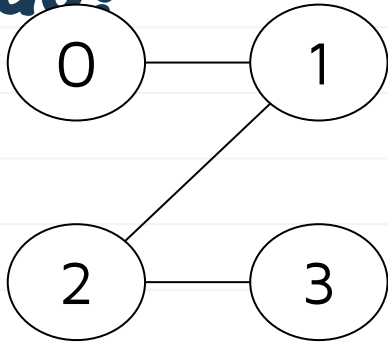
Exercise 2: What does the following code do?



```
adj = [
    [1],
    [0, 2],
    [1, 3],
    [2]
]
```

```
for u in range(len(adj)):
    for v in adj[u]:
        print(f"({u}, {v})")
```

Exercise 2: What does the following code do?

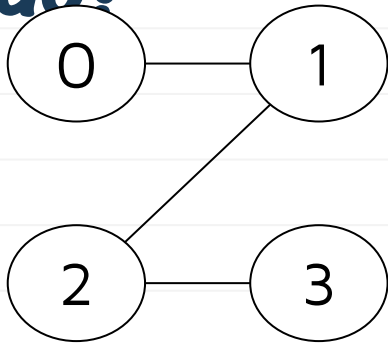


```
adj = [
    [1],
    [0, 2],
    [1, 3],
    [2]
]
```

```
for u in range(len(adj)):
    for v in adj[u]:
        print(f"({u}, {v})")
```

Prints all edges

Exercise 2: What does the following code do?



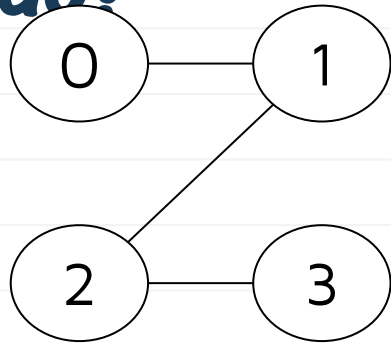
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    [1],
    [0, 2],
    [1, 3],
    [2]
]
```

```
for u in range(len(adj)):
    for v in adj[u]:
        print(f"({u}, {v})")
```

Prints all edges

(0, 1)

Exercise 2: What does the following code do?



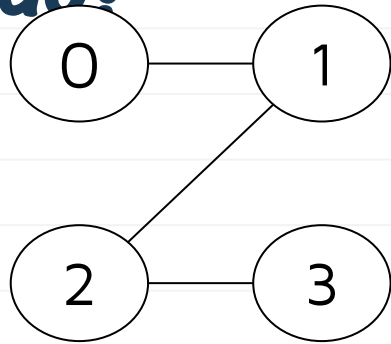
```
adj = [
    [1],
    [0, 2],
    [1, 3],
    [2]
]
```

```
for u in range(len(adj)):
    for v in adj[u]:
        print(f"({u}, {v})")
```

Prints all edges

(0, 1)
(1, 0)
(1, 2)

Exercise 2: What does the following code do?



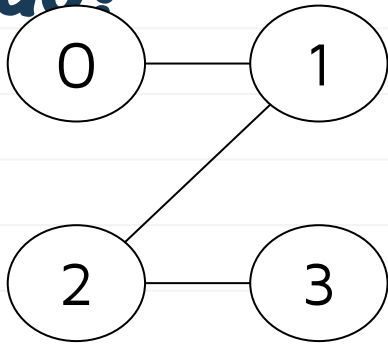
```
adj = [  
  [1],  
  [0, 2],  
  [1, 3],  
  [2]  
]
```

```
for u in range(len(adj)):  
    for v in adj[u]:  
        print(f"({u}, {v})")
```

Prints all edges

(0, 1)
(1, 0)
(1, 2)
(2, 1)
(2, 3)

Exercise 2: What does the following code do?



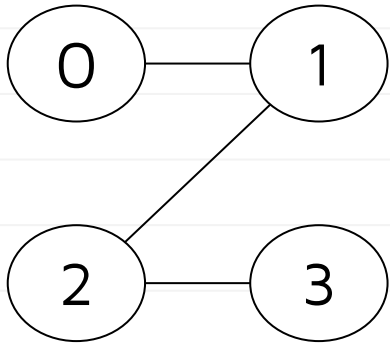
```
adj = [  
  [1],  
  [0, 2],  
  [1, 3],  
  [2]  
]
```

```
for u in range(len(adj)):  
    for v in adj[u]:  
        print(f"({u}, {v})")
```

Prints all edges

(0, 1)
(1, 0)
(1, 2)
(2, 1)
(2, 3)
(3, 2)

Exercise 3: Complexity



```
adj = [
    [1],
    [0, 2],
    [1, 3],
    [2]
]
```

```
for u in range(len(adj)):
    for v in adj[u]:
        print(f"({u}, {v})")
```

How many times is the **print** statement executed in terms of $|V|$ and $|E|$ (if undirected)?

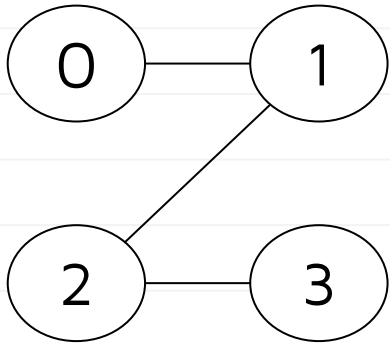
A: $|E| + |V|$

C: $2 * |E|$

B: $|E| * |V|$

D: $|E| + |V| - 1$

Exercise 3: Complexity



```
adj = [
    [1],
    [0, 2],
    [1, 3],
    [2]
]
```

```
for u in range(len(adj)):
    for v in adj[u]:
        print(f"({u}, {v})")
```

How many times is the **print** statement executed in terms of $|V|$ and $|E|$ (if undirected)?

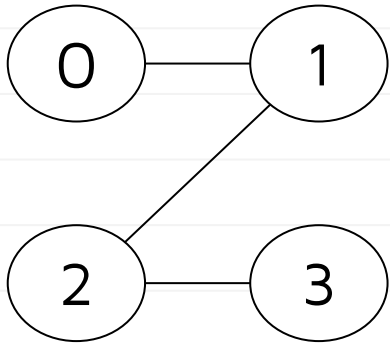
A: $|E| + |V|$

C: $2 * |E|$

B: $|E| * |V|$

D: $|E| + |V| - 1$

Exercise 3: Complexity



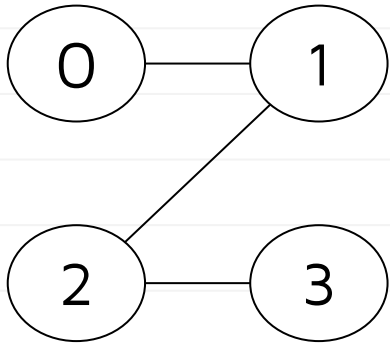
```
adj = [
    [1],
    [0, 2],
    [1, 3],
    [2]
]
```

```
for u in range(len(adj)):
    for v in adj[u]:
        print(f"({u}, {v})")
```

How many times is the **print** statement executed in terms of $|V|$ and $|E|$ (if directed)?

A: $|E|$

Exercise 4: Complexity



adj = [
[1],
[0, 2],
[1, 3],
[2]
]

```
for u in range(len(adj)):  
    for v in adj[u]:  
        print(f"({u}, {v})")
```

What is the time complexity in terms of $|V|$ and $|E|$?

A: $\Theta(|E| + |V|)$ C: $\Theta(|E|)$

B: $\Theta(|E| * |V|)$ D: $\Theta(|V|)$

Exercise 5:

```
adj = [  
    [],  
    [],  
    [],  
    []  
]
```

```
for u in range(len(adj)):  
    for v in adj[u]:  
        print(f"({u}, {v})")
```

0

1

2

3

What is the time complexity for this case?

A: $\Theta(|E| + |V|)$

C: $\Theta(|E|)$

B: $\Theta(|E| * |V|)$

D: $\Theta(|V|)$

Exercise 5:

```
adj = [  
    [],  
    [],  
    [],  
    []  
]
```

```
for u in range(len(adj)):  
    for v in adj[u]:  
        print(f"({u}, {v})")
```

0

1

2

3

What is the time complexity for this case?

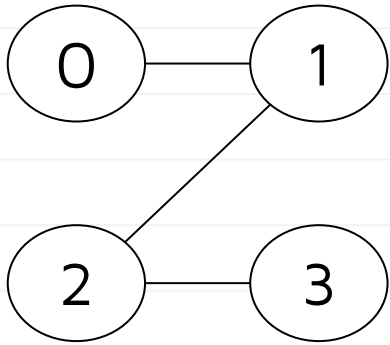
A: $\Theta(|E| + |V|)$

C: $\Theta(|E|)$

B: $\Theta(|E| * |V|)$

D: $\Theta(|V|)$

Exercise 4: Complexity



adj = [
[1],
[0, 2],
[1, 3],
[2]
]

```
for u in range(len(adj)):  
    for v in adj[u]:  
        print(f"({u}, {v})")
```

What is the time complexity in terms of $|V|$ and $|E|$?

A: $\Theta(|E| + |V|)$ C: $\Theta(|E|)$

B: $\Theta(|E| * |V|)$ D: $\Theta(|V|)$

Looping Over Edges

- Looping over all edges in this way takes $\Theta(V + E)$ time.
- In aggregate, the print statement is executed:
 - $2|E|$ times if graph is undirected.
 - $|E|$ times if graph is directed.
- This is called an **aggregate analysis**.

Space Requirements

- Need $\Theta(|V|)$ space for **outer list**.
- Plus $\Theta(|E|)$ space for **inner lists**.
- **In total:** $\Theta(|V| + |E|)$ space.

Example: Facebook

- Facebook has 2 billion users, 400 billion friendships.
- If each edge requires 32 bits: $(2 \text{ bits} \times 200 \times (2 \text{ billion}))$
 - = $64 \times 400 \times 10^9$ bits
 - = 3.2 terabytes
 - = 0.04 years of HD video

Dictionary of Sets

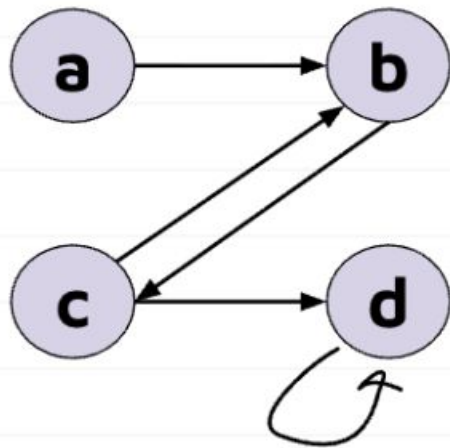
Trade Offs

- **Adjacency matrix**: fast edge query, lots of space.
- **Adjacency list**: slower edge query, space efficient.
- Can we have the best of both?

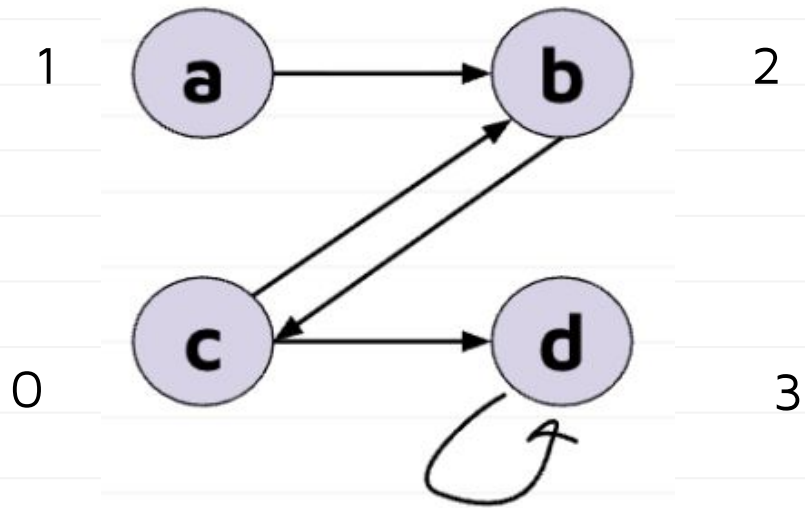
Idea

- Use hash tables.
- Replace inner edge lists by **sets**.
- Replace outer list with **dict**.
 - Doesn't speed things up, but allows nodes to have arbitrary labels.

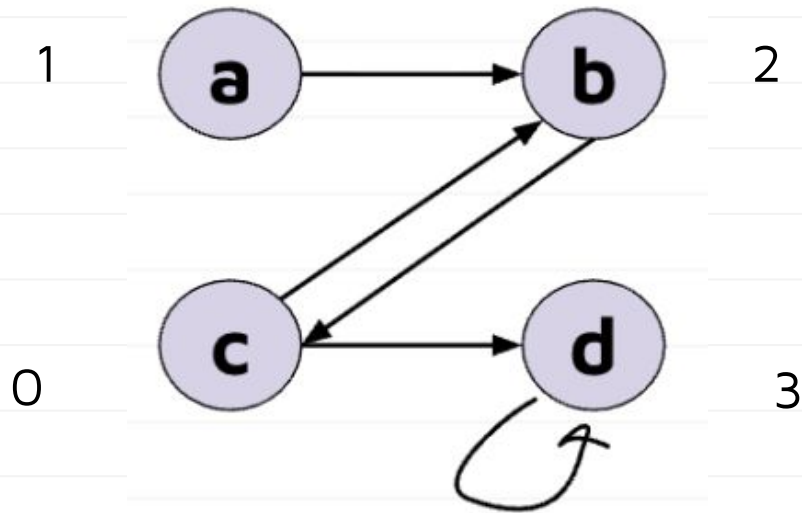
Example



Example



Example



```
adj = {
```

```
    'a':
```

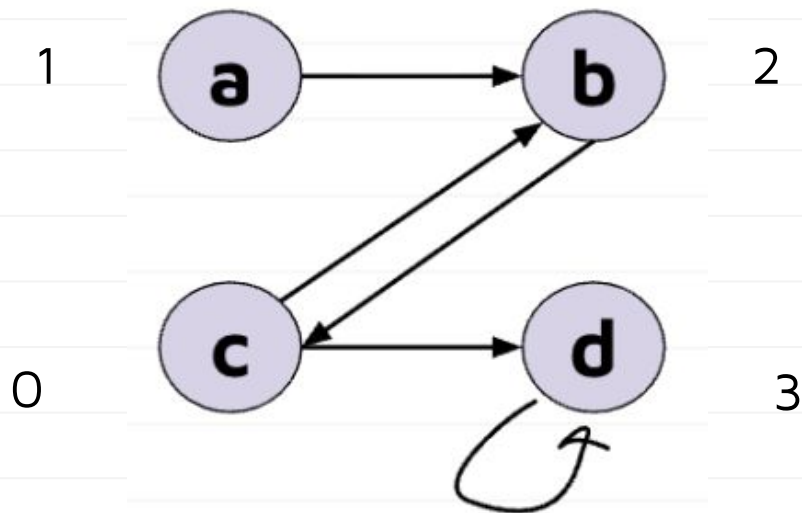
```
    'b':
```

```
    'c':
```

```
    'd':
```

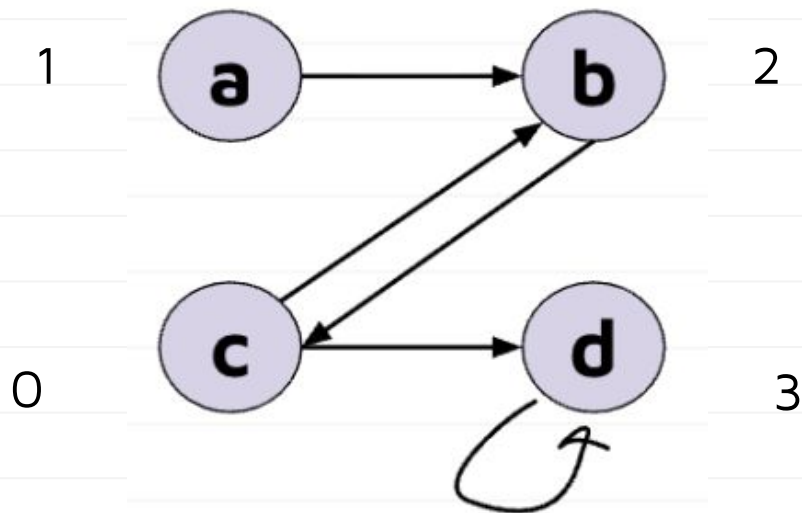
```
}
```

Example



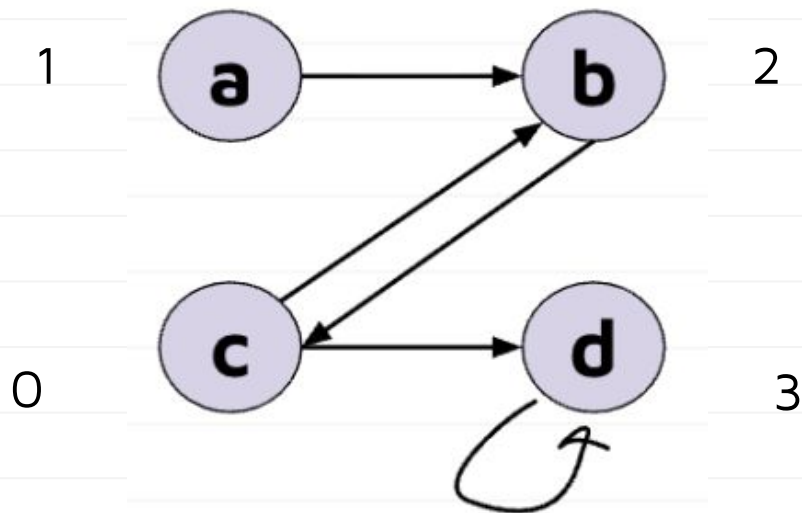
```
adj = {  
    'a': {'b'},  
    'b':  
    'c':  
    'd':  
}
```

Example



```
adj = {  
    'a': {'b'},  
    'b': {'c'},  
    'c':  
    'd':  
}
```

Example



```
adj = {
```

```
    'a': {'b'},
```

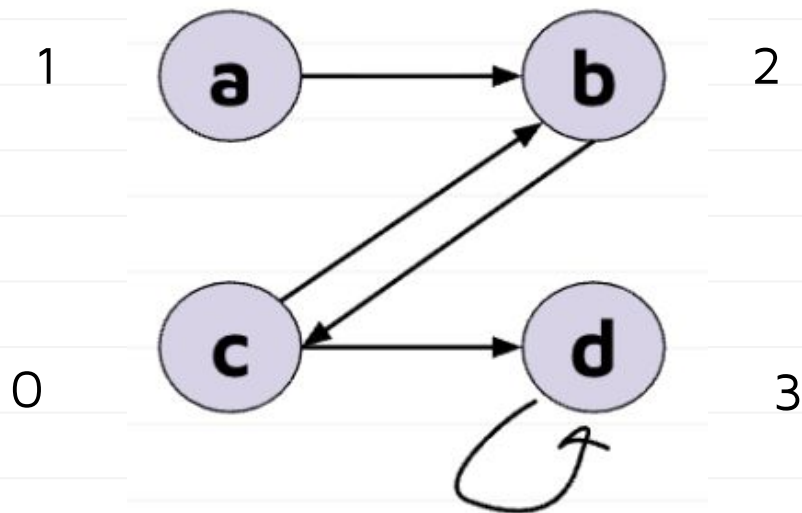
```
    'b': {'c'},
```

```
    'c': {'b', 'd'},
```

```
    'd':
```

```
}
```

Example



adj = {

'a': {'b'},

'b': {'c'},

'c': {'b', 'd'},

'd': {'d'}

}

Time Complexity

operation	code	time
Edge query	<code>j in adj[i, j]</code>	$\Theta(1)$ on average
Degree(i)	<code>len(adj[i])</code>	$\Theta(1)$ on average

Space Requirements

- Requires only $\Theta(V + E)$.
- But there is overhead to using hash tables.

Dict-of-sets implementation

- Install with `pip install dsc40graph`
- Import with `import dsc40graph`
- **Docs:** <https://eldridgejm.github.io/dsc40graph/>
- **Source code:** <https://github.com/eldridgejm/dsc40graph>
- Will be used in HW coding problems.



Thank you!

Do you have any questions?

CampusWire!