

DSC 40B

***Lecture 5 : Best,
Worst Cases***

Correction (mic)

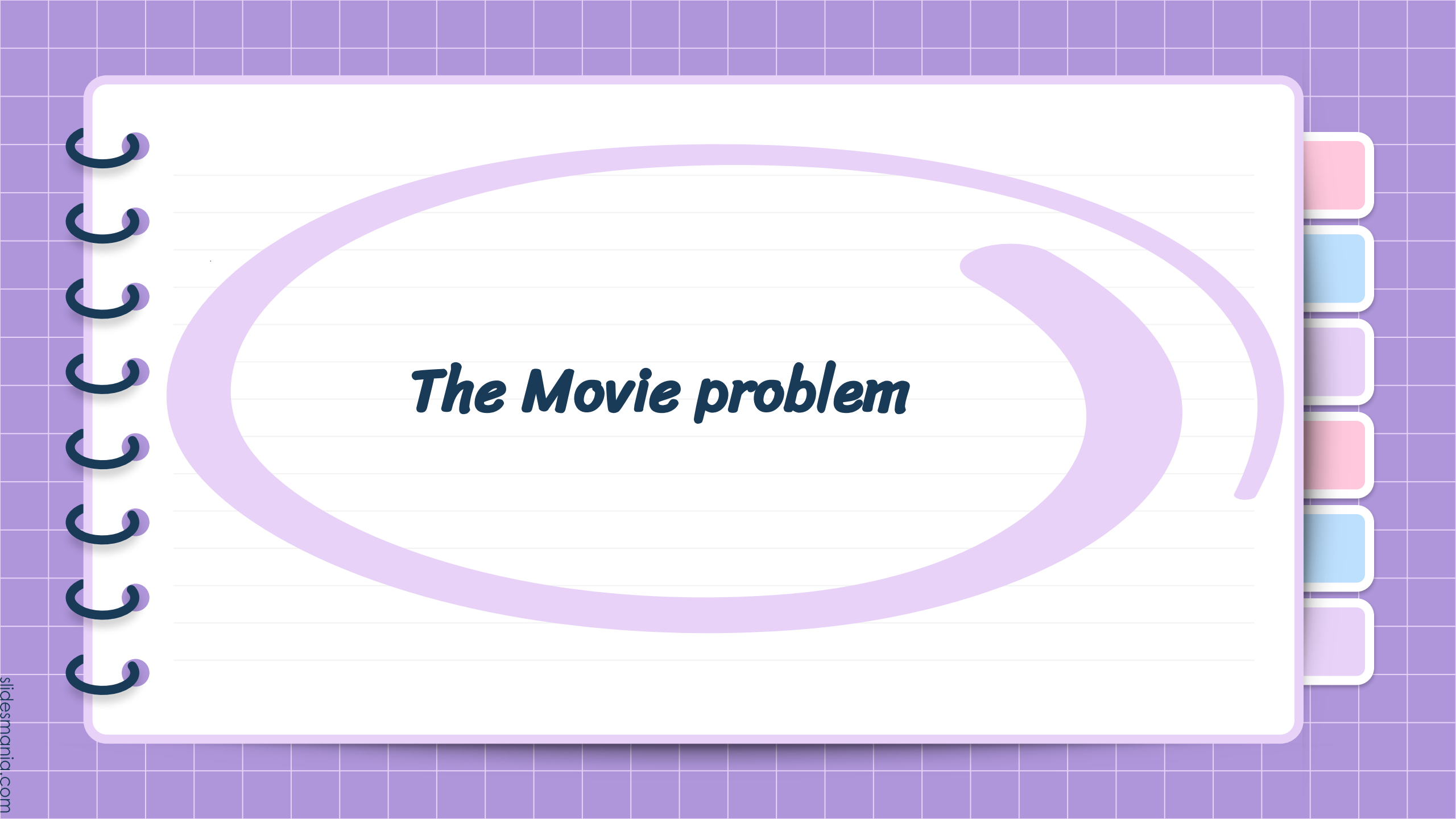
https://docs.google.com/presentation/d/1JsgA03S0rMCrWD6oORdvjVB1cdt0HQHVHEx6zszylkA/edit?slide=id.g388f030f327_0_12#slide=id.g388f030f327_0_12



Agenda

Plan for the 5-6 lectures

- Best, Worst and Average cases.



The Movie problem

The Movie Problem



The Movie Problem

- **Given:** an array `movies` of movie durations, and the flight duration `t`
- **Find:** two movies whose durations add to `t`
 - If no two movies sum to `t`, return **None**.

Exercise

- Design a brute force solution to the problem. What is its time complexity?

Brute force

```
def find_movies(movies, t):  
    n = len(movies)  
    for i in range(n):  
        for j in range(i + 1, n):  
            if movies[i] + movies[j] == t:  
                return (i, j)  
    return None
```

Time Complexity

- It looks like there is a **best** case and **worst** case.
- How do we formalize this?

For the future...

- Can you come up with a better algorithm?
- What is the *best possible* time complexity?



Best and Worst Cases

Definition

- Define $T_{\text{best}}(n)$ to be the **least** time taken by the algorithm on **any input of size n** .
- The asymptotic growth ($n \rightarrow \text{infinity}$) of $T_{\text{best}}(n)$ is the algorithm's **best case time complexity**.

Example 1: mean. Best case?

```
def mean(arr):  
    total = 0  
    for x in arr:  
        total += x  
    return total / len(arr)
```

Example 1: mean

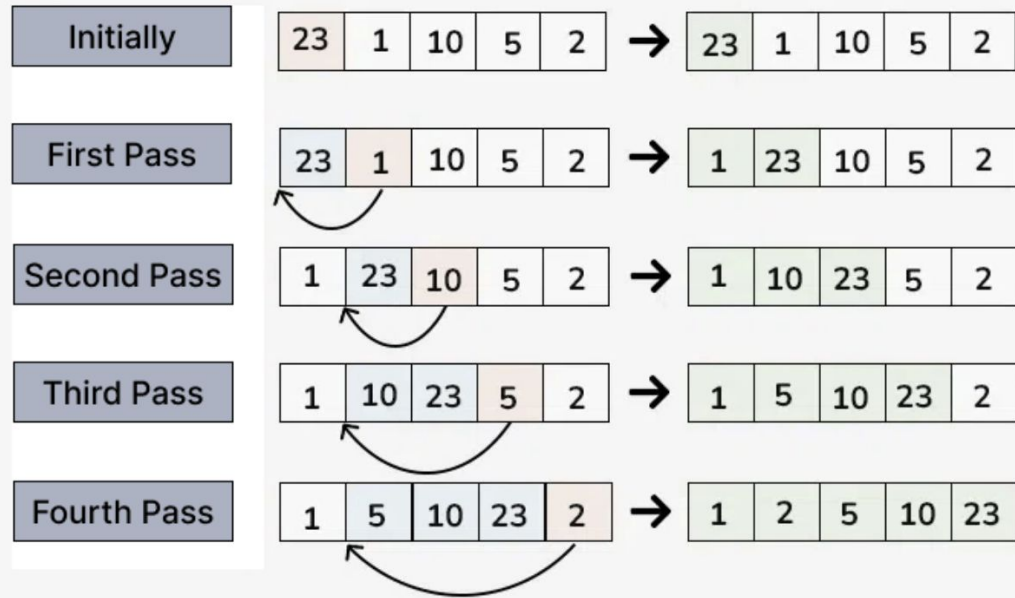
```
def mean(arr):  
    total = 0  
    for x in arr:  
        total += x  
    return total / len(arr)
```

$$T_{\text{best}}(n) = n$$

Caution!

- The best case is **never**: “the input is of size one or empty”.
- The best case is about the **structure** of the input, not its **size**.
- Not always constant time!
 - **Example**: sorting.

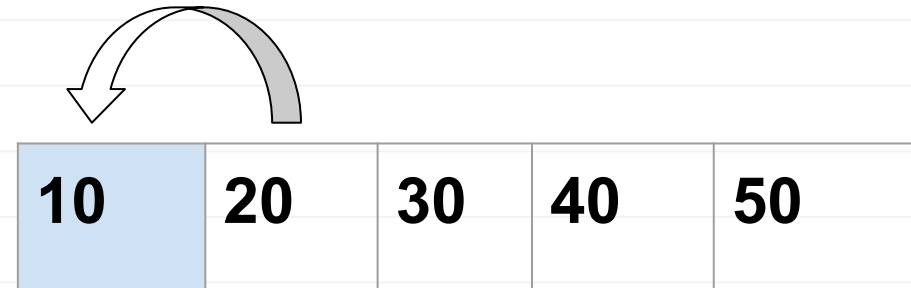
```
def insertionSort(arr):  
    for i in range(1, len(arr)):  
        key = arr[i]  
        j = i - 1  
        while j >= 0 and key < arr[j]:  
            arr[j + 1] = arr[j]  
            j -= 1  
        arr[j + 1] = key
```



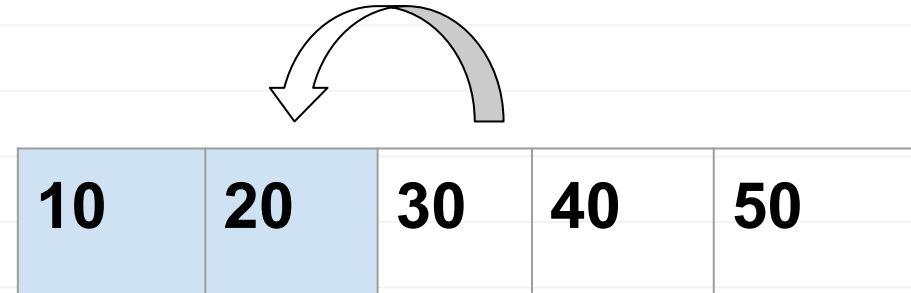
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```

10	20	30	40	50
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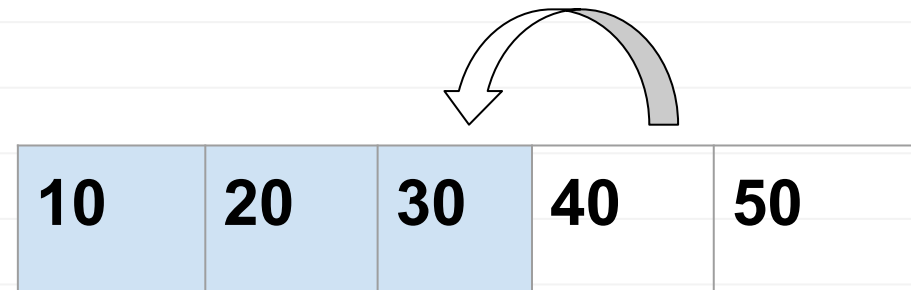
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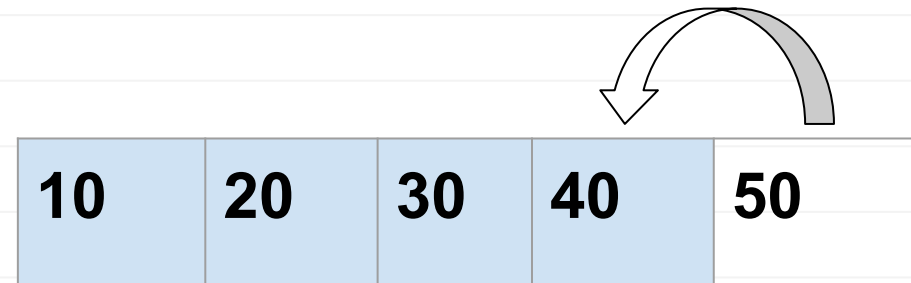
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        arr[j + 1] = key
```

$$T_{\text{best}}(n) = n$$

10	20	30	40	50
----	----	----	----	----

Time Complexity of mean

- Linear time **for all inputs**, $\Theta(n)$.
- Depends **only** on the array's **size**, n , not on its actual elements.

Example 2: Linear Search

- **Given:** an array `arr` of numbers and a target `t`.
- **Find:** the index of `t` in `arr`, or `None` if it is missing.
- **Example:** `arr = [-3, -6, 7, 3, 0, 15, 4]`
 - `t` is 7
 - 2 is returned (index of 7)

Cool function!

```
def linear_search(arr, t):  
    for i, x in enumerate(arr):  
        if x == t:  
            return i  
    return None
```

Exercise: Time complexity? Best?

```
def linear_search(arr, t):  
    for i, x in enumerate(arr):  
        if x == t:  
            return i  
    return None
```

A: Constant

B: n

C: Something else

Exercise: Time complexity? Worst?

```
def linear_search(arr, t):  
    for i, x in enumerate(arr):  
        if x == t:  
            return i  
    return None
```

Observation

- It looks like there are two extreme cases...

*The **Best** Case*

- When the target, t , is the very first element.
- The loop exits after one iteration.
- $\Theta(1)$ time?

*The **Worst** Case*

- When the target, t , is **not** in the array at all.
- The loop exits after n iterations.
- $\Theta(n)$ time?

Time Complexity

- `linear_search` can take vastly different amounts of time on two inputs of the **same size**.
 - Depends on **actual elements** as well as size.
- It has no single, overall time complexity.
- Instead we'll report **best** and **worst** case time complexities.

Best Case Time Complexity

- How does the time taken in the **best case** grow as the input gets larger?

Best Case

- In linear_search's **best case**, $T_{\text{best}}(n) = c$, no matter how large the array is.
- The **best case time complexity** is $\Theta(1)$.

Worst Case Time Complexity

- How does the time taken in the **worst case** grow as the input gets larger?

Definition

- Define $T_{\text{worst}}(n)$ to be the **most time** taken by the algorithm on any input of size n .
- The asymptotic growth of $T_{\text{worst}}(n)$ is the algorithm's **worst case time complexity**.

Worst Case

- In the worst case, `linear_search` iterates through the entire array.
- The **worst case** time complexity is $\Theta(n)$.

Exercise: times: Best and Worst

```
def func(arr):  
    n = len(arr)  
    for x in arr:  
        for y in arr:  
            if x + y == 10  
                return sum(arr)
```

A: $\Theta(1)$

B: $\Theta(n)$

C: $\Theta(n^2)$

D: $\Theta(n^3)$

Best Case

- Best case occurs when the first element is 5: $5 + 5 = 10$
- `sum(arr)` takes $\Theta(n)$ time
- Exists, taking $\Theta(n)$ time in total

Worst Case

- Worst case occurs when no two numbers add to 10.
- Has to loop over all $\Theta(n^2)$ pairs.
- Worst case time complexity: $\Theta(n^2)$.
- **Note:** Not $\Theta(n^3)$ since the `sum(arr)` only runs once.

Note

- An algorithm like `linear_search` **doesn't** have one single time complexity.
- An algorithm like `mean` **does**, since the best and worst case time complexities coincide.

Main Idea

Reporting **best** and **worst** case time complexities gives us a richer of the performance of the algorithm.



Thank you!

Do you have any questions?

CampusWire!