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**DSC 40B - Homework 07**  
 Due: Tue 11:59pm (PT) — NOT Wed 11:59pm

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**No slip days allowed (solutions released Wed).**

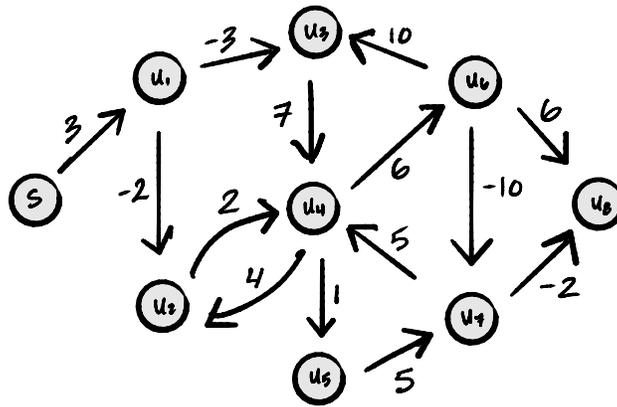
Write your solutions to the following problems by either typing them up or handwriting them on another piece of paper. Unless otherwise noted by the problem's instructions, show your work or provide some justification for your answer. Homeworks are due via Gradescope at 11:59 p.m.

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**Problem 1.**

Run Bellman-Ford on the following graph using node  $s$  as the source. Below each node  $u$ , write the shortest path length from  $s$  to  $u$ . Mark the predecessor of  $u$  by highlighting it or making a bold arrow. You can assume that `graph.edges` produces the graph's edges in the following order:

$(u_2, u_4), (u_7, u_8), (u_6, u_3), (s, u_1), (u_1, u_2), (u_4, u_5), (u_1, u_3),$   
 $(u_3, u_4), (u_6, u_7), (u_4, u_6), (u_5, u_7), (u_7, u_4), (u_4, u_2), (u_6, u_8)$

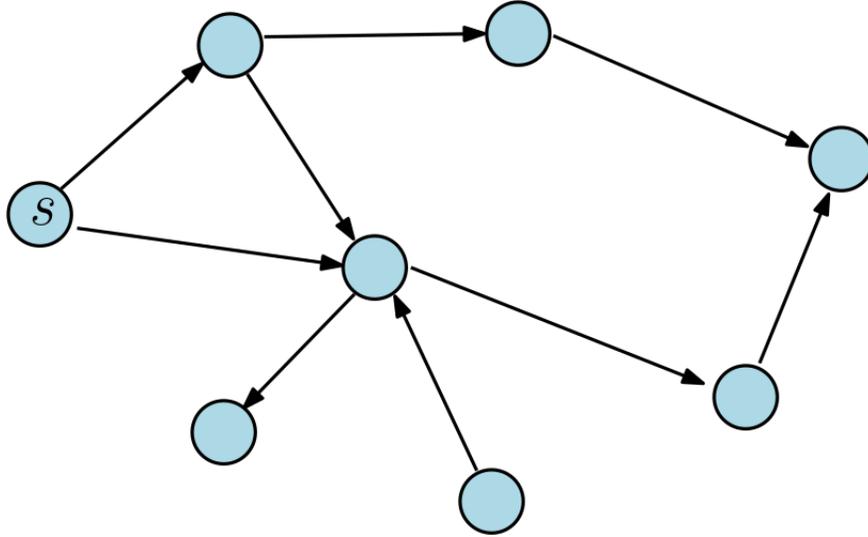


**Problem 2.**

Recall that the Bellman-Ford algorithm (with early stopping) will terminate early if, after updating every edge, no predecessors have changed. Suppose it is known that the greatest shortest path distance in a graph  $G = (V, E)$  has  $\Theta(\sqrt{V})$  edges. What is the worst case time complexity of Bellman-Ford when run on this graph? State your answer using  $\Theta$  notation.

**Problem 3.**

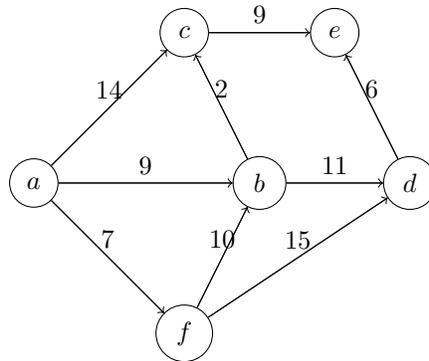
Suppose we are given a directed and weighted graph  $G = (V, E)$  that look like the following figure. It is known that there are no negative cycles.



Suppose the Bellman-Ford algorithm **with early stopping** is run on this graph shown above using node  $s$  as the source (the edge weights are purposely not shown). In the worst case, how many iterations of the outer for-loop of Bellman-Ford will be performed? Briefly justify your answer.

**Problem 4.**

Run Dijkstra's Algorithm on the following graph using node  $a$  as the source. Below each node  $u$ , write the shortest path length from  $a$  to  $u$ . Mark the predecessor of  $u$  by highlighting it or making a bold arrow.



**Problem 5.**

True or False. Suppose  $G = (V, E, \omega)$  is a weighted graph for which all edges are positive, except for those edges of a node  $s$  which may or may not be negative. If Dijkstra's algorithm is run on  $G$  with  $s$  as the source, the correct shortest paths will be found. Assume that the graph does not have any negative loops. Justify your answer.

**Problem 6.**

Suppose  $G = (V, E, \omega)$  is a weighted graph without negative cycles. Suppose we have an edge  $(u, v) \in E$  whose weight is  $w(u, v) = 5$ . Let  $s \in V$  be the source node. Let  $\delta(s, u)$  and  $\delta(s, v)$  denote the shortest path distance from  $s$  to  $u$  and to  $v$ , respectively. For each statement below, just answer True or False. No justification or examples are required.

- a) (True or False): Now suppose that  $G$  is a **directed** graph. Then, it is possible that  $\delta(s, u) = 50$  but  $\delta(s, v) = 10$ .
- b) (True or False): Now suppose that  $G$  is a **undirected** graph. Then, it is possible that  $\delta(s, u) = 50$  but  $\delta(s, v) = 10$ .

**Problem 7.**

Consider the code for Bellman-Ford (without early stopping) which was given in lecture:

```
def bellman_ford(graph, weights, source):
    est = {node: float('inf') for node in graph.nodes}
    est[source] = 0
    predecessor = {node: None for node in graph.nodes}

    for i in range(len(graph.nodes) - 1):
        for (u, v) in graph.edges:
            update(u, v, weights, est, predecessor)

    return est, predecessor
```

Note that the algorithm assumes there is a specific ordering of edges in the graph and for each iteration of the outer-for loop, the inner for-loop will update all edges in this order. Now imagine the Bellman-Ford algorithm is run on the graph below (edge weights not marked) using  $a$  as the source node.



Please provide an **ordered list of** all edges, such that if we run the Bellman-Ford algorithm on this ordering of edges, then the distance estimate at node  $d$  **will not** be correct **until** after the outer-for loop finishes three iterations.