DSC 40B - Discussion 01

Problem 1.

What is the time complexity of the following functions? State your answer using Θ notation.

Solution: $\Theta(n^3)$

```
b) def foo(n):
    while n > 1:
        n /= 10
        print(n)
```

Solution: $\Theta(\log n)$

```
c) def foo(n):
    for i in range(n):
        for j in range(i**2): # <-- notice the bound!
            print(i + j)</pre>
```

```
Solution: \Theta(n^3)
```

```
d) def pairs(numbers):
    result = []
    for x in numbers:
        for y in numbers:
            result.append((x, y))
```

return result

Solution: $\Theta(n^2)$

e) def foo(numbers):

for pair in pairs(numbers):
 print(sum(pair))

Solution: $\Theta(n^2)$

Note that the result of pairs(numbers) is actually only computed once, on the first iteration. On this first iteration, Python will try to produce the 0th element of pairs(numbers), which it will need to compute in $\Theta(n^2)$ time. After this result is computed, subsequent executions of the for loop line will take $\Theta(1)$ time as they simply produce the next element of the precomputed result. So, this function is equivalent to:

```
def foo(numbers):

lst = pairs(numbers) # \theta(n^2)

for pair in lst: # \theta(n^2)

print(sum(pair)) # \theta(1)
```

Problem 2.

Let $f(n) = \sum_{p=0}^{n} 3^{p}$. What is f in Θ notation?

Solution:

General form of a geometric sum $\sum_{p=0}^{n} x^p = \frac{1-x^{n+1}}{1-x}$. Substituting our equation yields $\sum_{p=0}^{n} 3^p = \frac{1-3^{n+1}}{1-3}$. Therefore, $f(n) = \Theta(3^n)$ after throwing out the constants.

Problem 3.

Consider the code below where heights is an array of n elements:

```
for i in range(n):
    for j in range(2*i):
        height = heights[i] + heights[j]
```

What is the time complexity of the code?

Solution:

We can see that outer iteration runs n times, for inner iteration:

On outer iter 1, inner body runs 0 times

On outer iter 2, inner body runs 2 times

On outer iter 3, inner body runs 4 times

Hence,

On outer iter α , inner body runs $(2\alpha - 2)$ times

$$\sum_{\alpha=1}^{n} f(\alpha) = \sum_{\alpha=1}^{n} 2 * \alpha - 2$$
$$\sum_{\alpha=1}^{n} 2 * \alpha - 2 = 0 + 2 + 4 + \dots + (2n - 4) + (2n - 2)$$

This is an arithmetic series and we know the formula for sum of an arithmetic series is:

$$Sum = \frac{n}{2} * (a_1 + a_n)$$

Where n is the number of terms, a_1 is the first term in the series and a_n is the last term. Therefore sum of the series:

$$Sum = \frac{n}{2} * (0 + 2n - 2) = n(n - 1) = n^{2} - n$$

Therefore the time complexity can be given as $\theta(n^2)$.